



IMT Atlantique

Bretagne-Pays de la Loire
École Mines-Télécom

INTRODUCTION TO INTERNET OF THINGS

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IMT ATLANTIQUE



Brest



Rennes



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Nantes



OCIF - OBJETS COMMUNICANTS ET INTERNET DU FUTUR



The Internet of Things

Mesh network vs long range radio

End-to-end connectivity

Consumer-producer loop

Security



Smart grid

Scalability



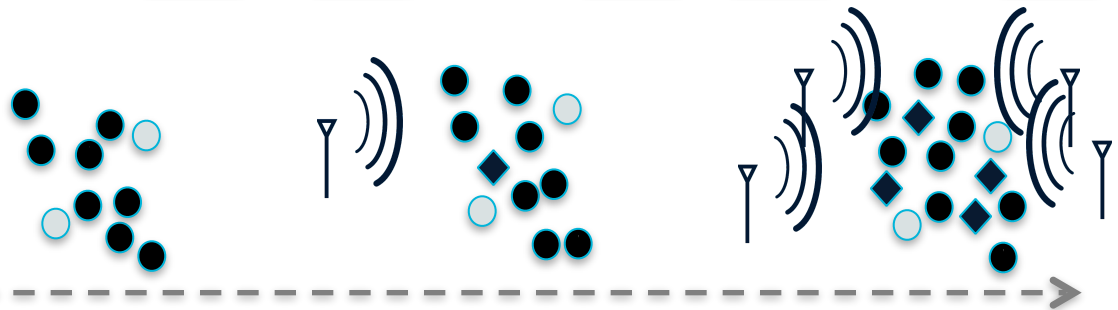
Smart city

Quality of service

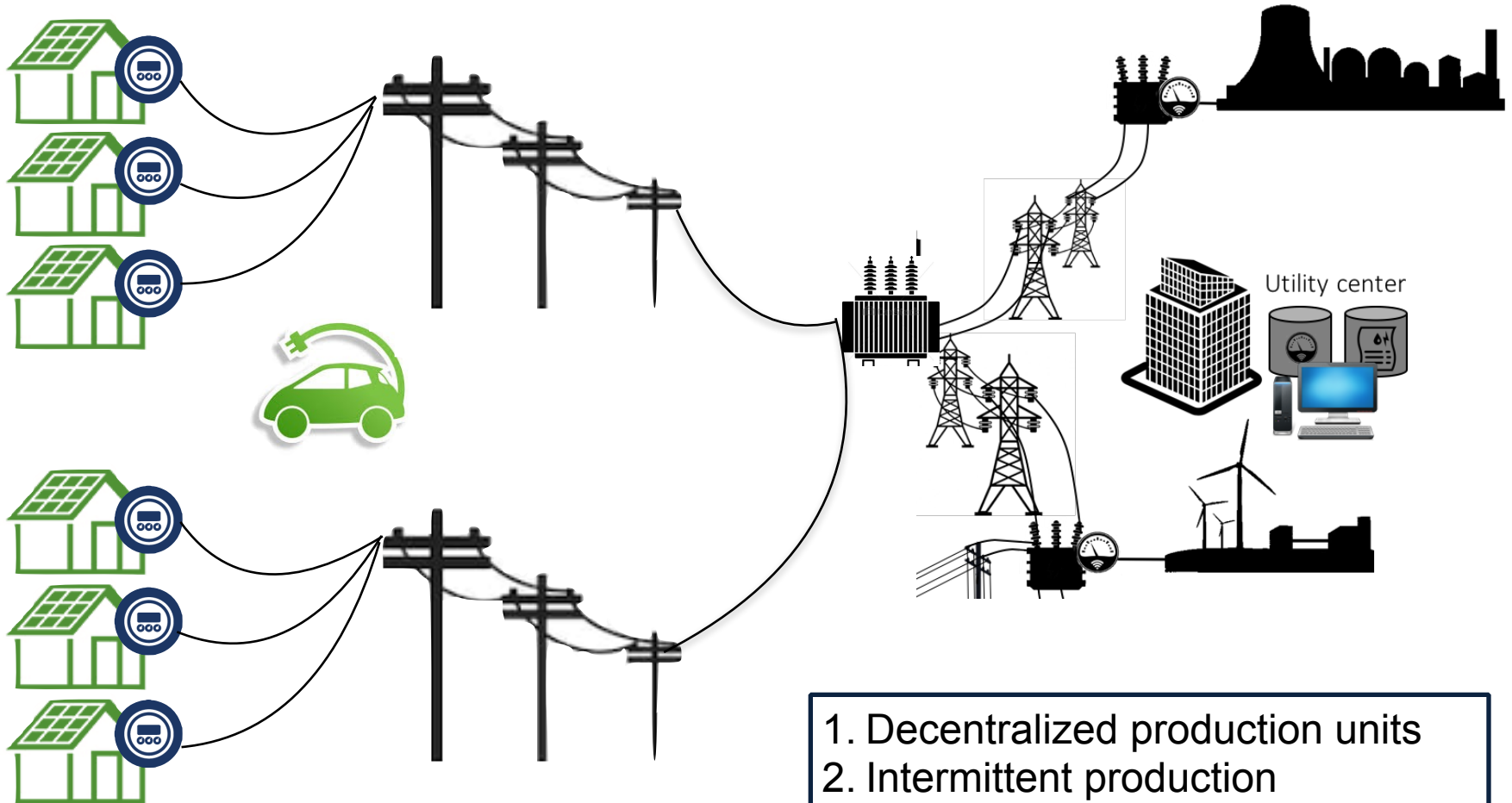


Industry 4.0

**Unified access network
for heterogeneous applications**

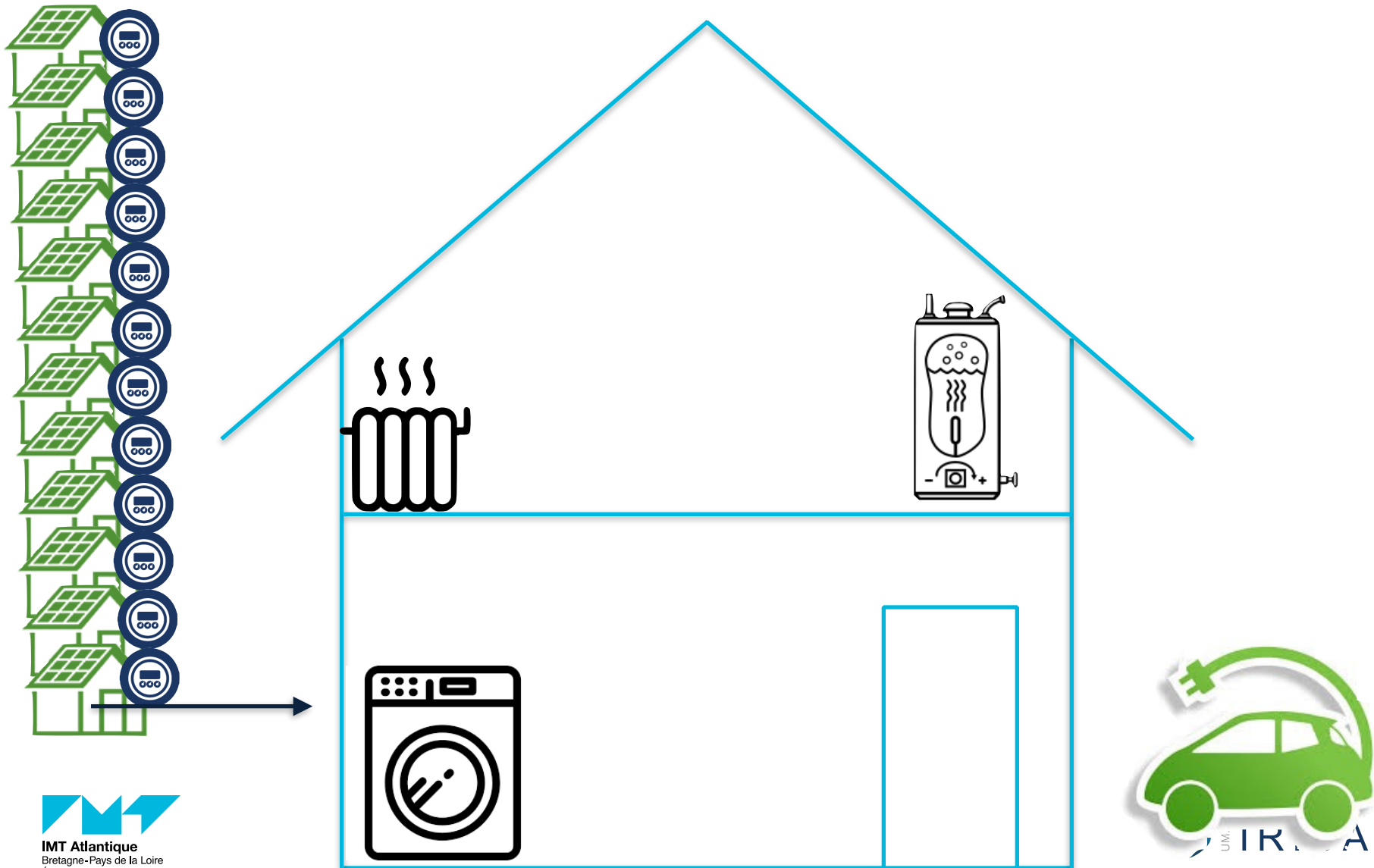


SMART GRID - HOW WE SEE THINGS



1. Decentralized production units
2. Intermittent production
 - ▶ Need for ICT system
 - ▶ Shift the demand
 - ▶ Monitor in real time

SMART GRID - HOW WE SEE THINGS



SMART GRID - SMART PLUG

Program

The Atmel Atmega128RFA1 microcontroller can easily run customized programs powered by the Contiki Operating System.

Monitor

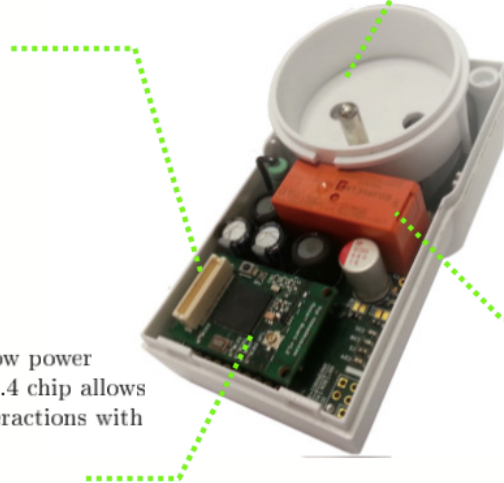
Plug in an electric device and retrieve informations about its electric consumption in real time.

Collect

The integrated low power transceiver 802.15.4 chip allows easy wireless interactions with one or more smartplugs.

Control

The Smartplug is equipped with a relay allowing the device the turn the current on or off.



INTRODUCTION

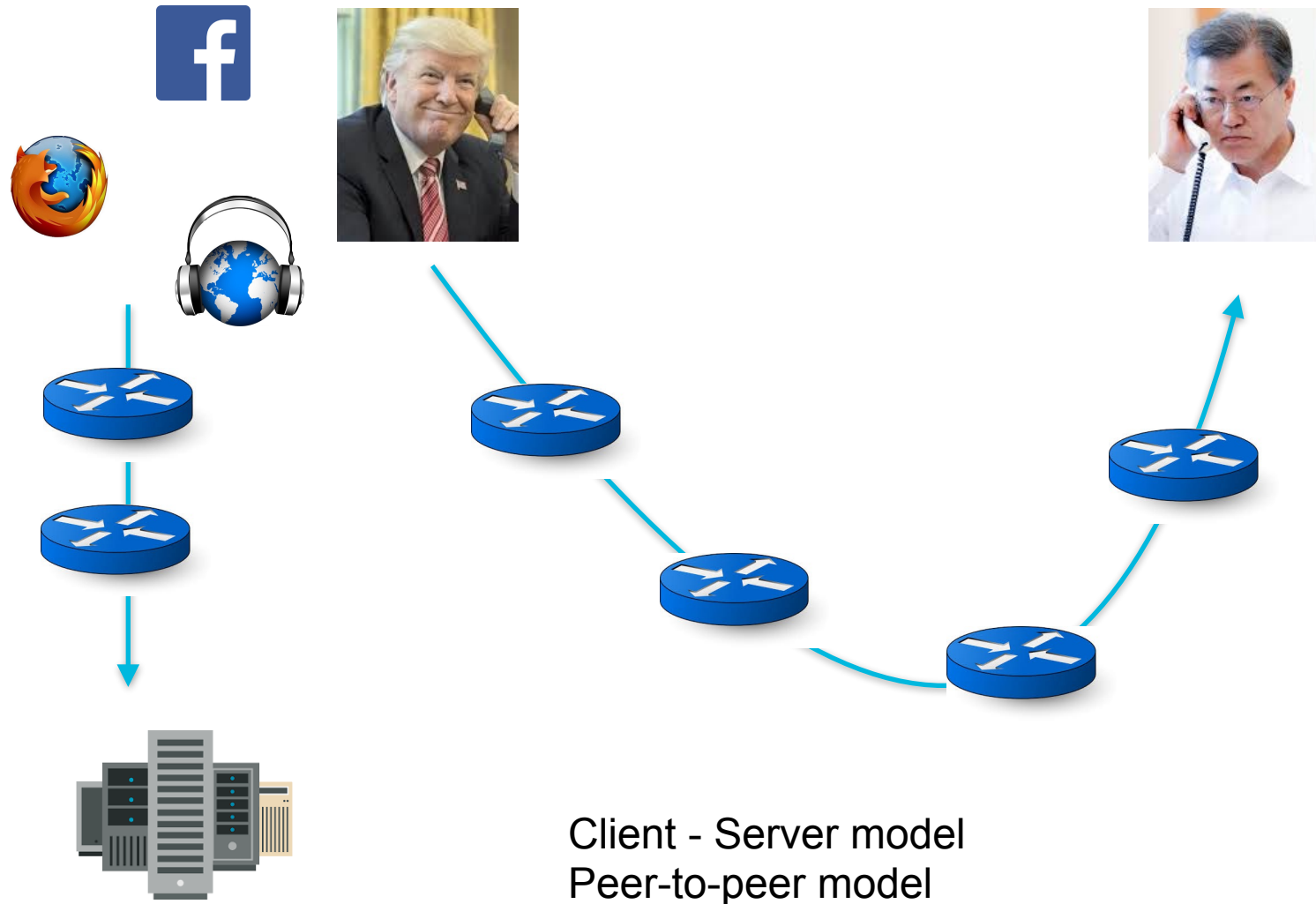


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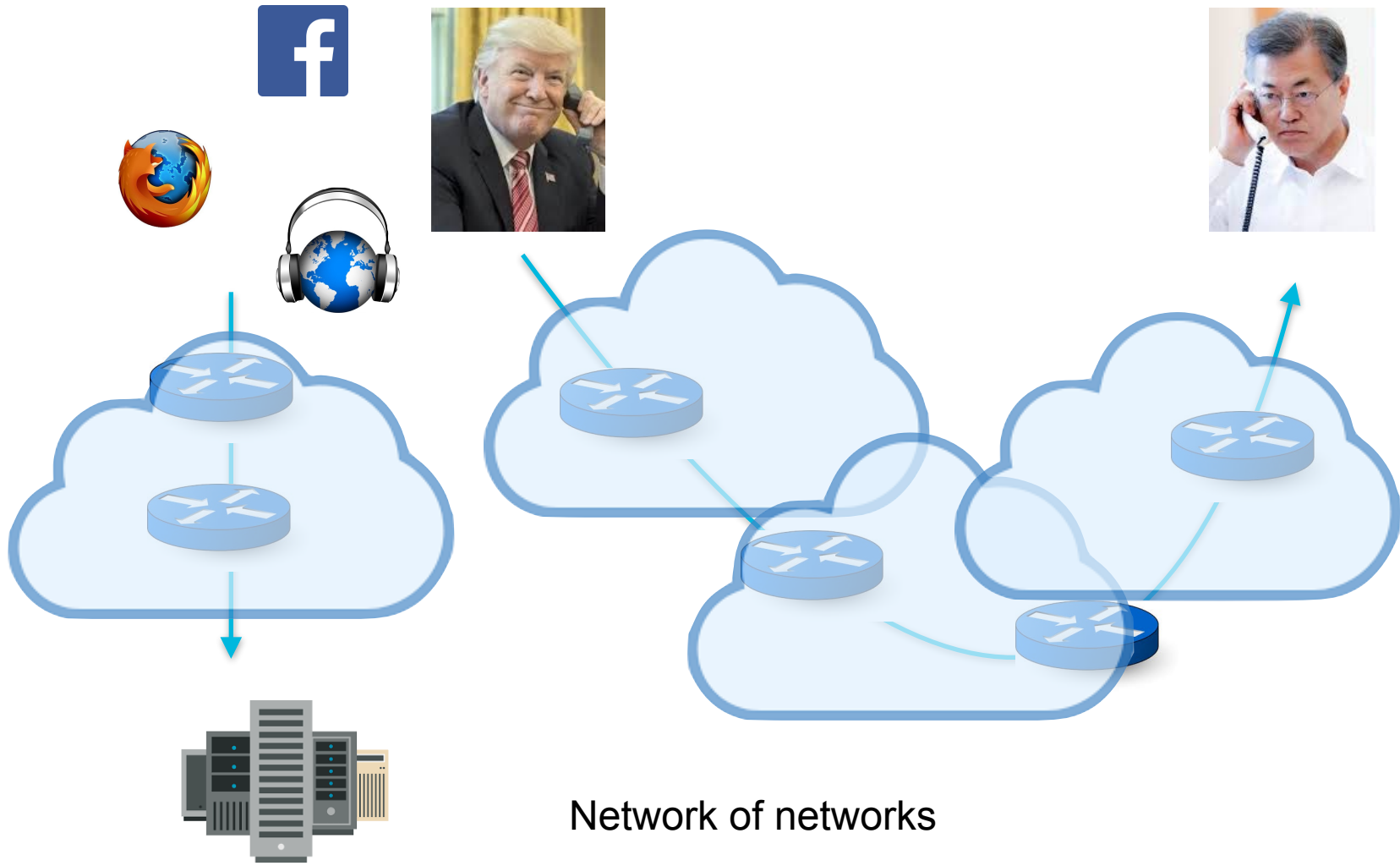
WHAT IS THE INTERNET?



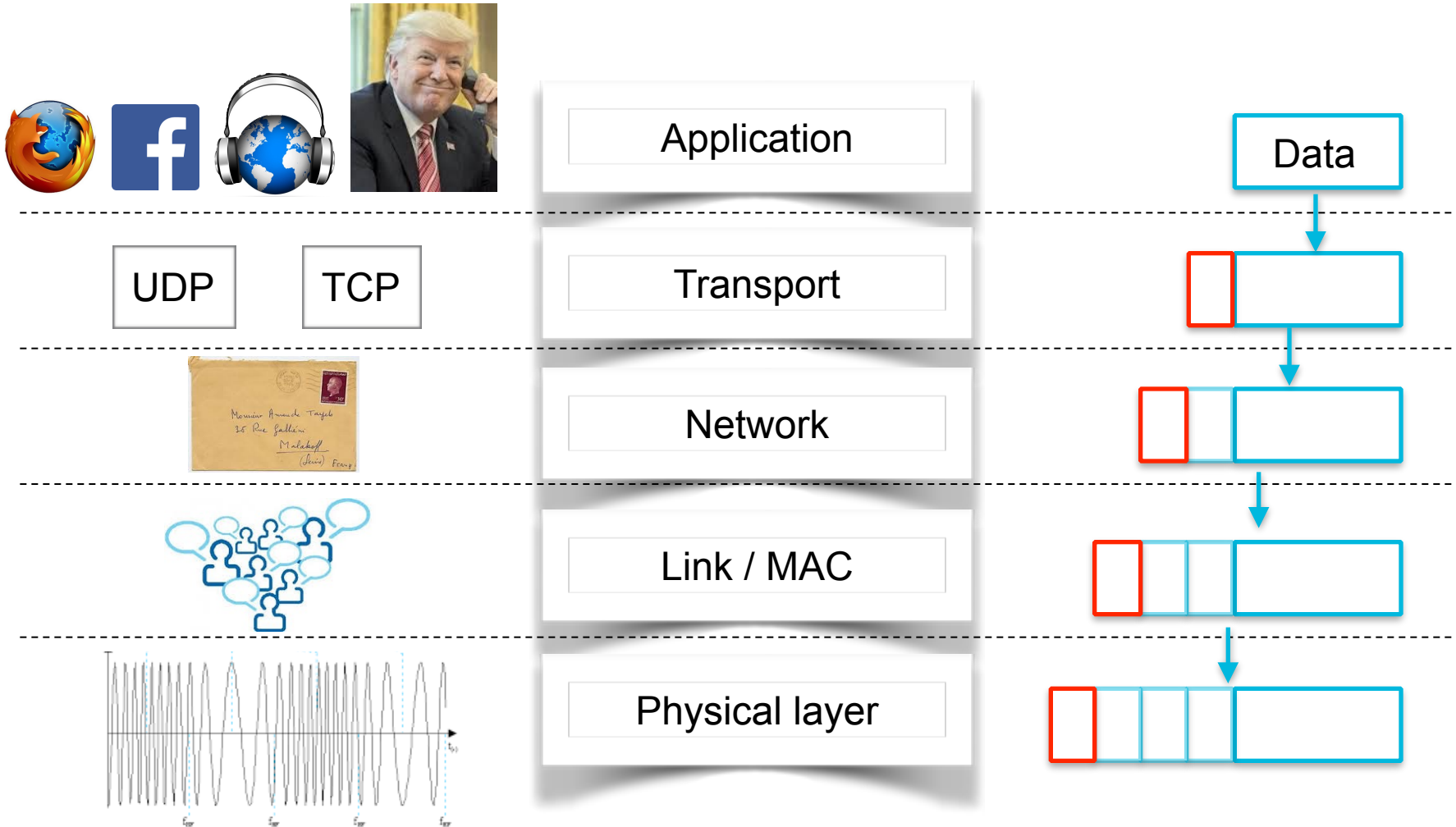
WHAT IS THE INTERNET?



WHAT IS THE INTERNET?



LAYERED APPROACH

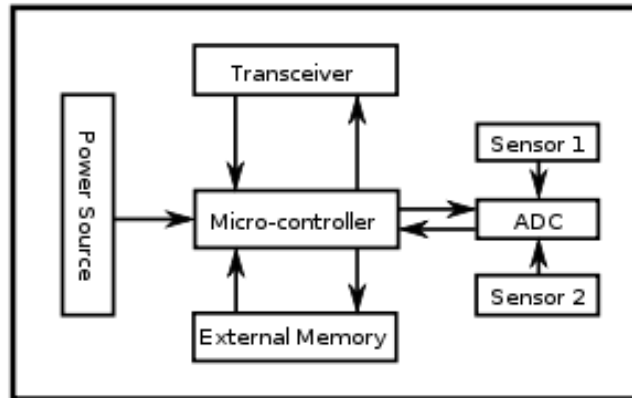


WHAT ARE THE STRENGTH OF THE INTERNET?

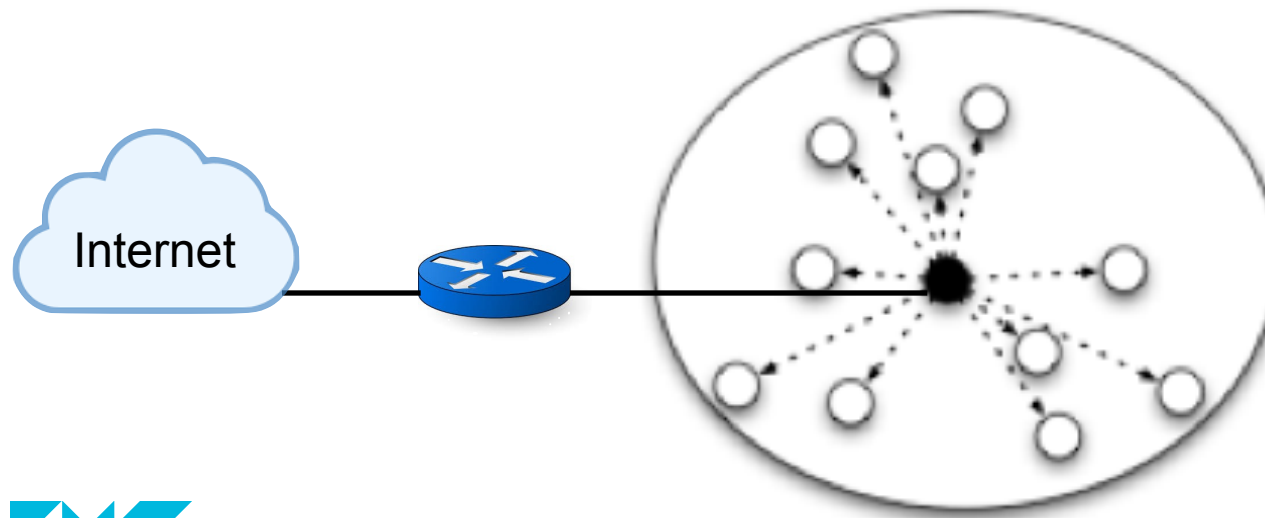
- Worldwide
- Standard
- Multi-applications
- Multi-technologies
- Works whatever the link capacity

So why not using it for the things (i.e., sensors, objects)?

THE INTERNET OF THINGS



Things



Connected

LAYERED APPROACH



Application

Transport

Network

Link / MAC

Physical layer



overhead data

Bluetooth
4.0
Low Energy



sigfox

LoRa

ISA
100
WIRELESS

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IRISA

LAYERED APPROACH



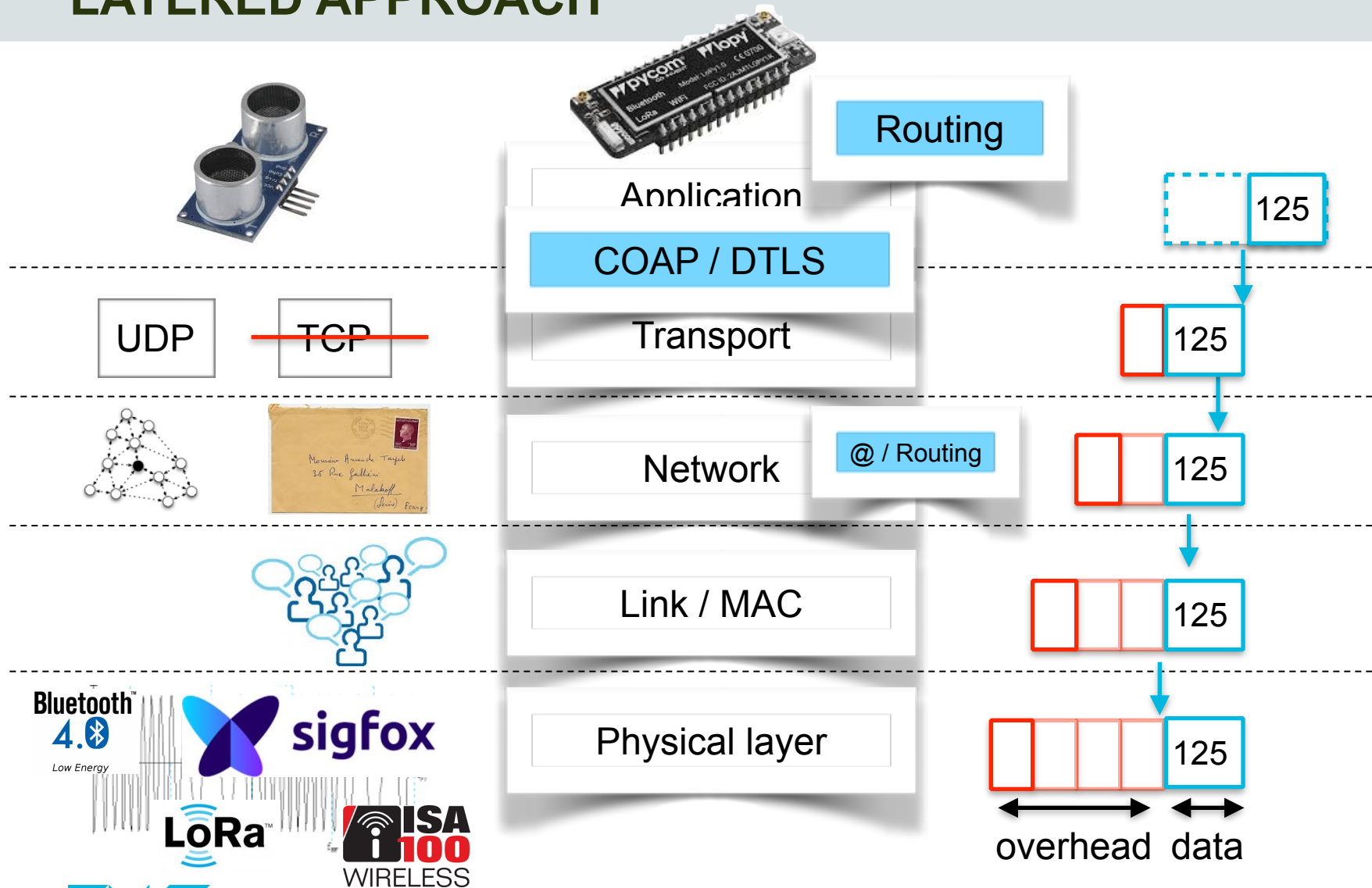
Long Range

	range	bitrate	frequency
SigFox	13 km	100 bps	900MHz
LTE-M	15 km	150 kbps – 1 Mbps	licensed 900MHz
LoRa	11 km	10 kbps	900MHz

Short Range

	range	bitrate	frequency
IEEE 802.15.4	10 m	250 Kbps	2.4GHz
Wi-Fi HaLow	1 km	150 kbps – 1 Mbps	900MHz
BLE	100 m	125 kbps – 2Mbps	2.4GHz

LAYERED APPROACH



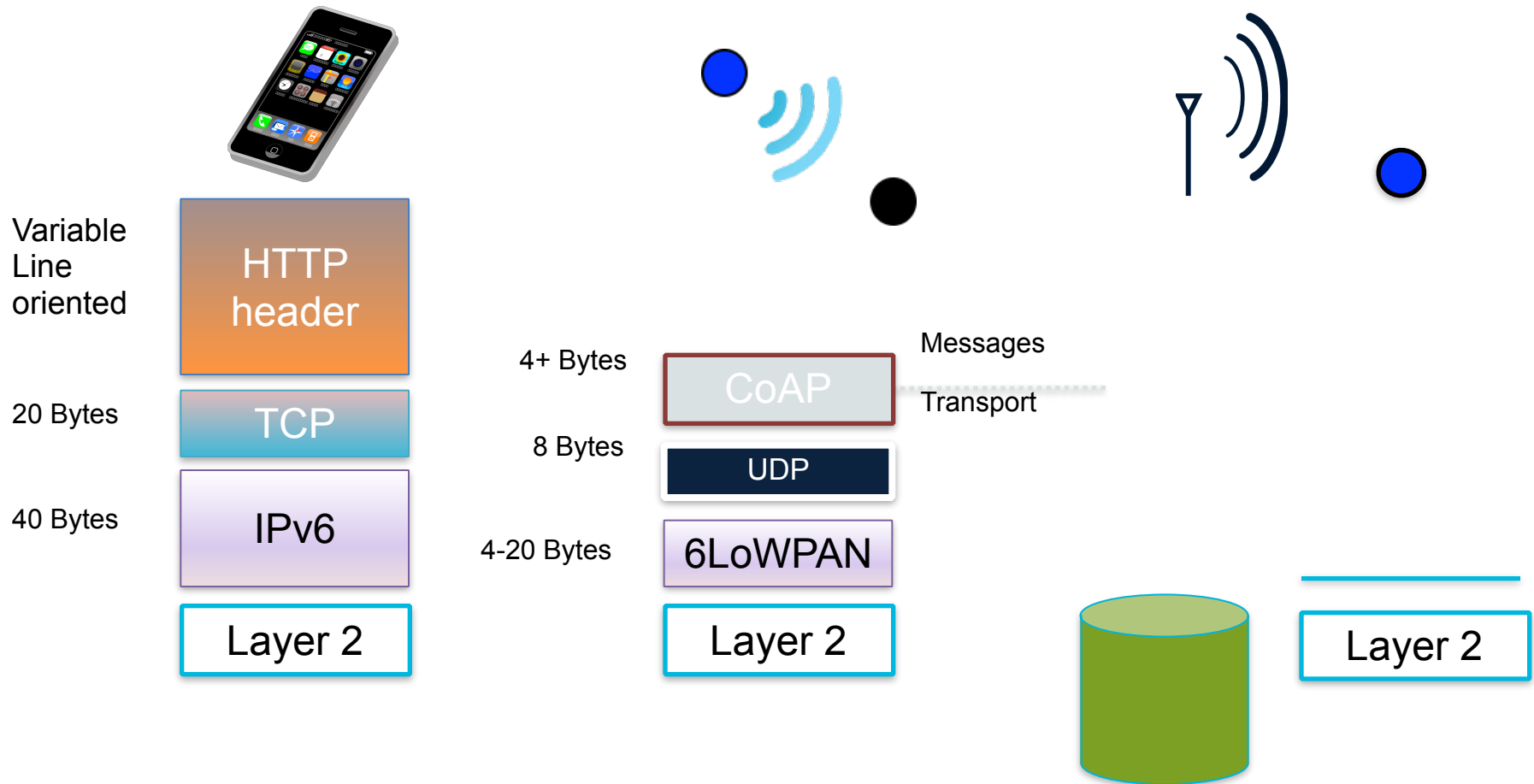
6LOWPAN

► Header Size Calculation. . .

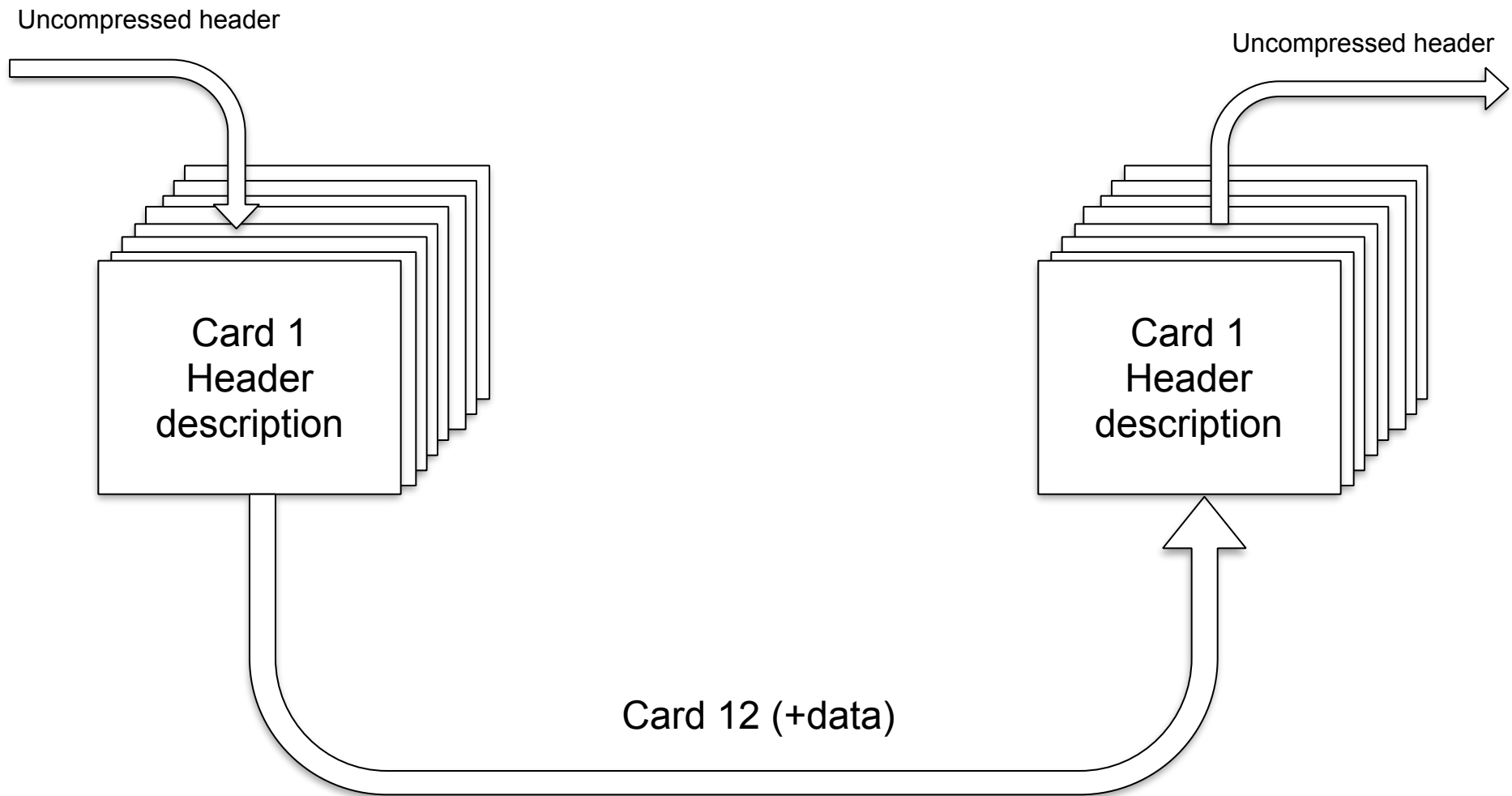
- IPv6 header is 40 bytes, UDP header is 8 bytes
- 802.15.4 MAC header can be up to 25 bytes (null security) or $25+21=46$ bytes (AES-CCM-128)
- With the 802.15.4 frame size of 127 bytes:
 $127-25-40-8 = 54$ bytes (null security)
 $127-46-40-8 = \mathbf{33 \text{ bytes}}$ (AES-CCM-128)

► IPv6 MTU Requirements

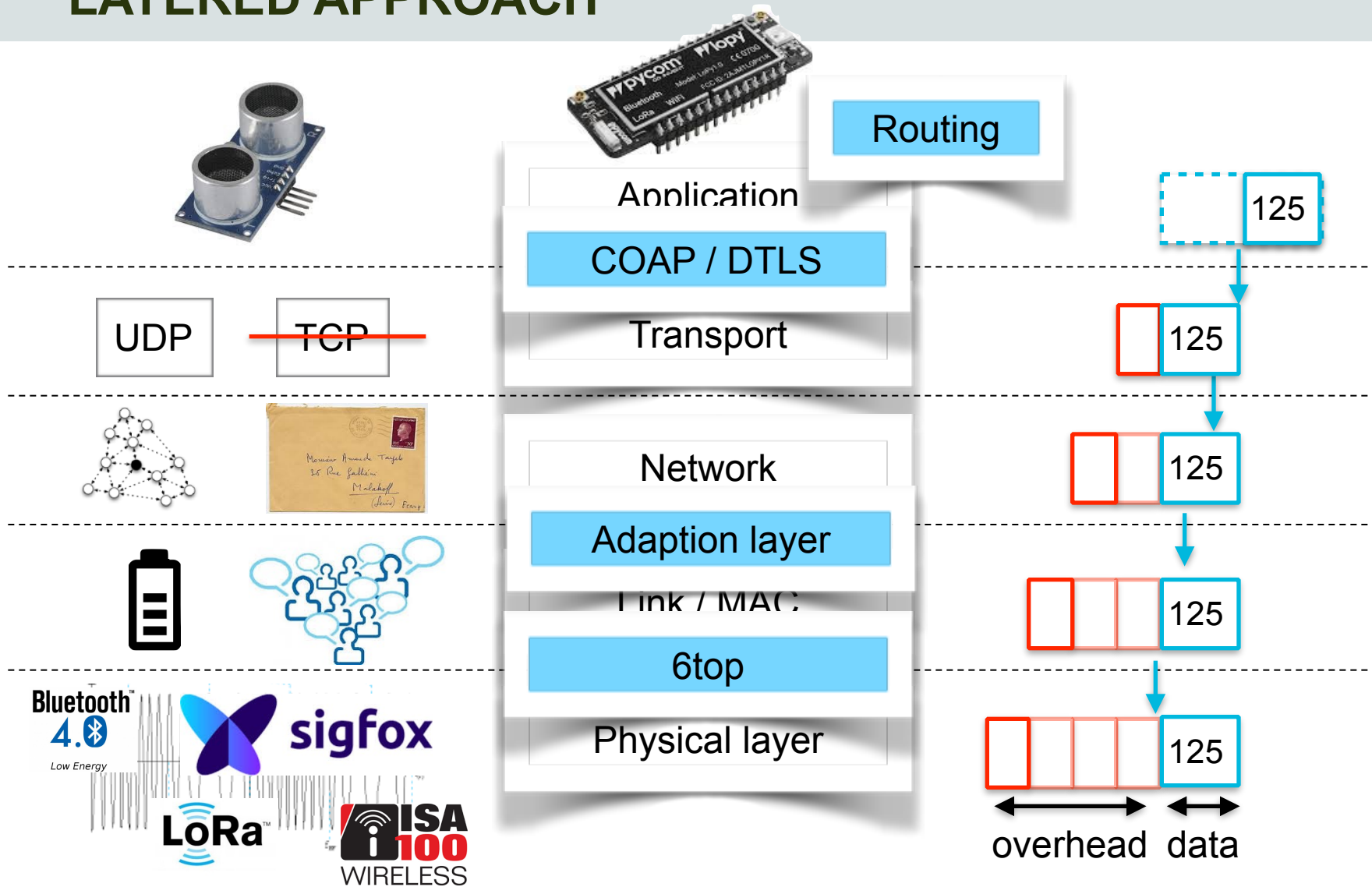
- IPv6 requires that links support an MTU of 1280 bytes
- Header compression &
- Link-layer fragmentation / reassembly is needed



Static Context Header Compression²⁰



LAYERED APPROACH



THE INTERNET OF THINGS

Benefits

- Worldwide
- Standard
- Multi-applications
- Multi-technologies
- Works whatever the link capacity

Challenges

- Adapted wireless technologies
 - Low energy consumption
 - Low capacity: low data rate, lossy, short frame, high delay
- Large overhead
 - Compression
- Data representation
- Routing

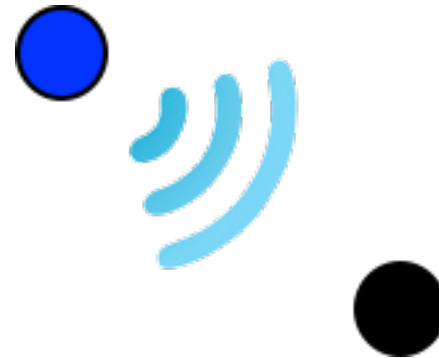
MEDIUM ACCESS CONTROL



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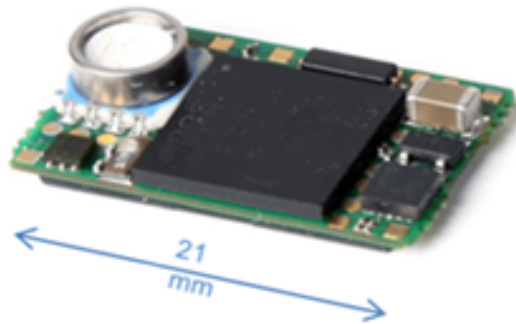
MEDIUM ACCESS CONTROL

- Objective: share the medium access
 - Who can send?
 - When a node should expect to receive data?
 - With the following goals
 - fairness
 - reliability
 - scalability
 - low latency
 - air throughput
 - **low energy consumption**



LET'S CONNECT THE THINGS

Sensor node

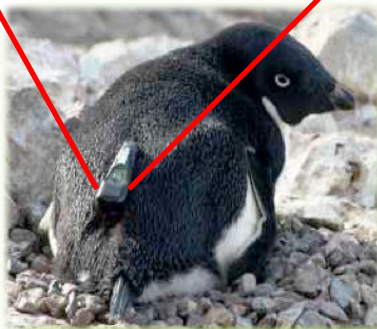


Battery lifetime: 2 years

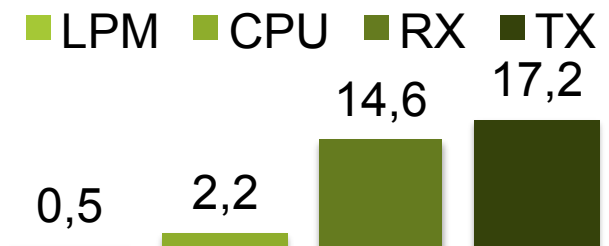
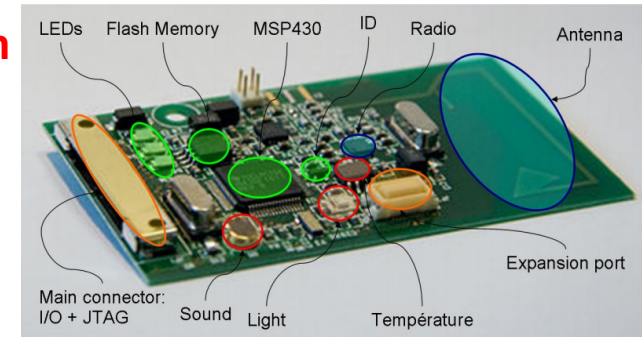
Battery lifetime: < 1 month



Bio-logger



Wireless sensor node



Electricity consumption (mA)
of electronic components [1]

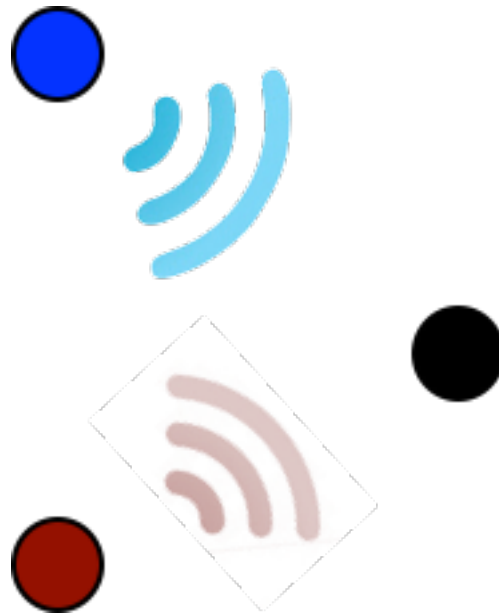


MEDIUM ACCESS CONTROL

- More than 100 MAC proposals for sensor networks
- Source of energy consumption
 - Receive / transmit
 - Idle listening
 - Control packets
 - Overhearings
 - Collisions

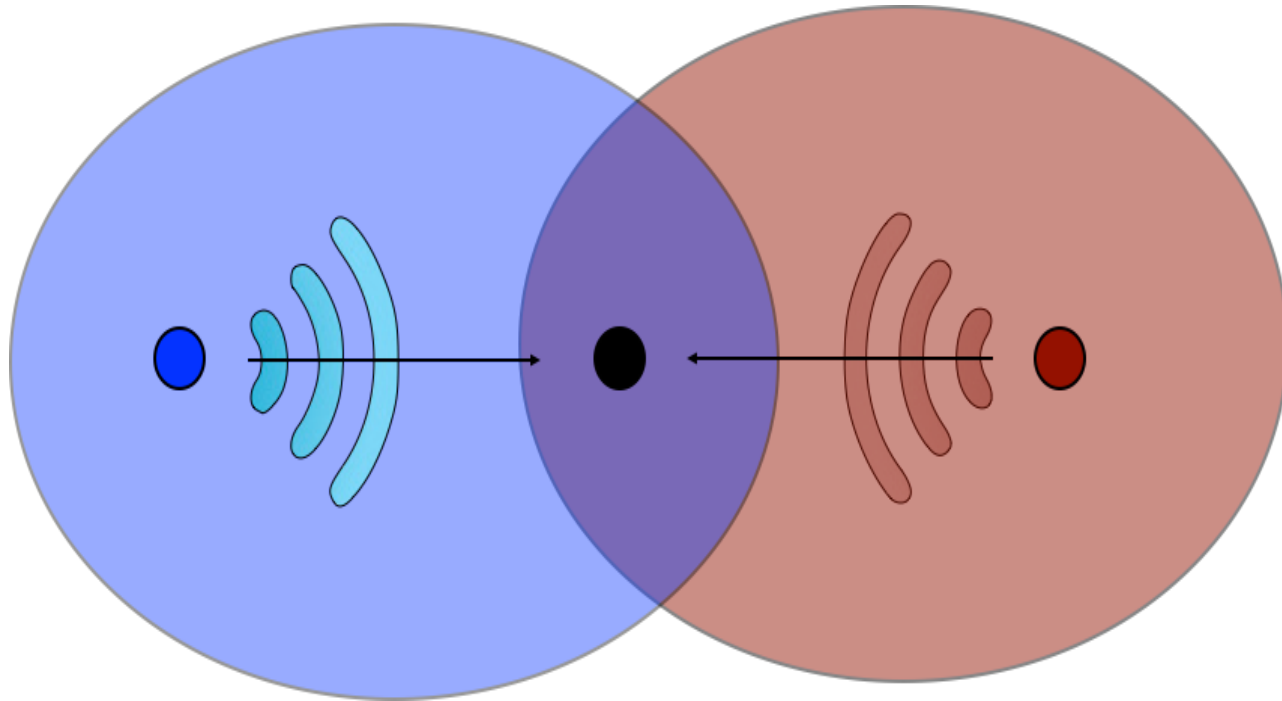
A COLLISION EXAMPLE

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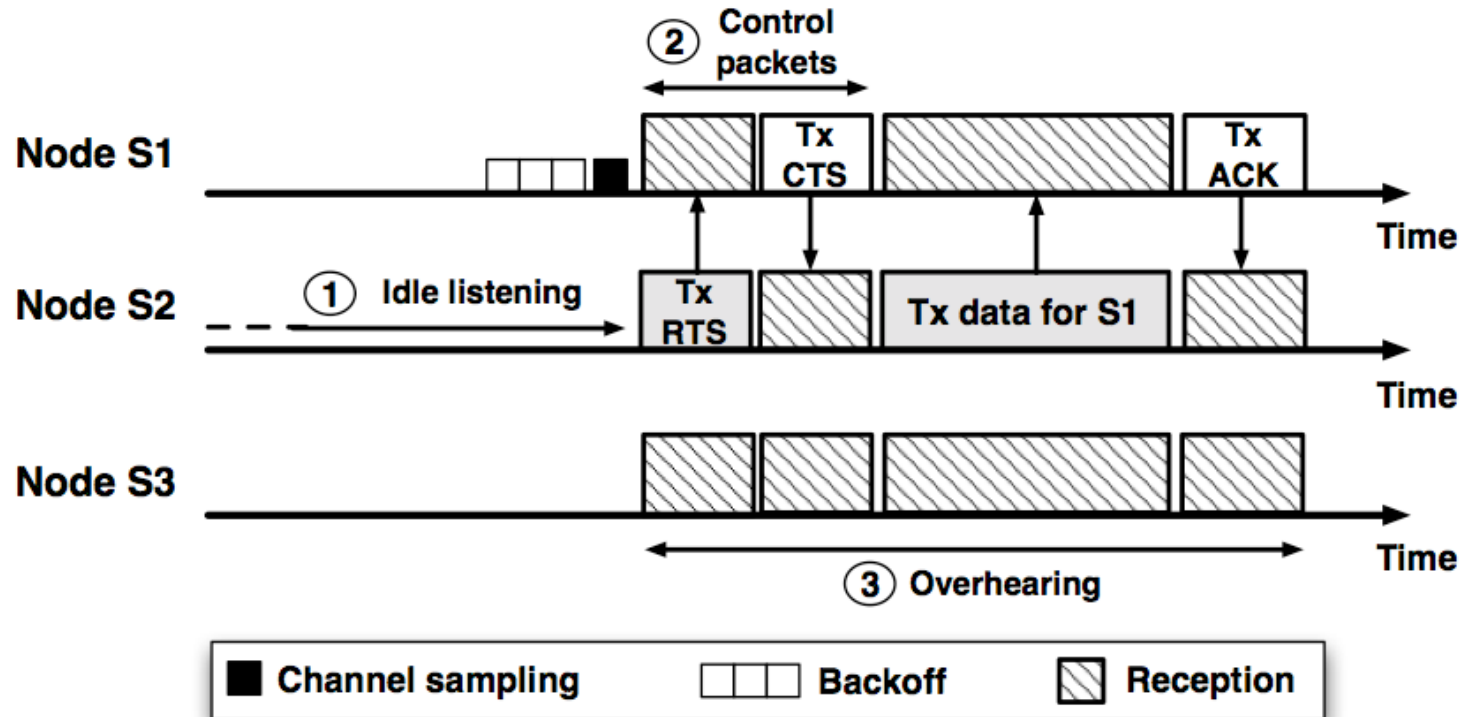


HIDDEN NODE PROBLEM

28

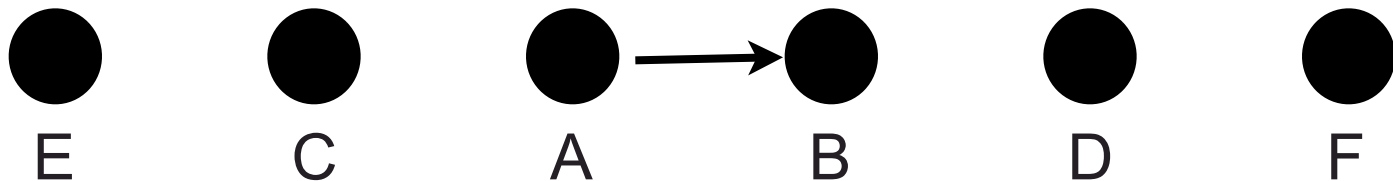


OVERHEARING AND IDLE LISTENING



ENERGY SAVING: SLEEP!

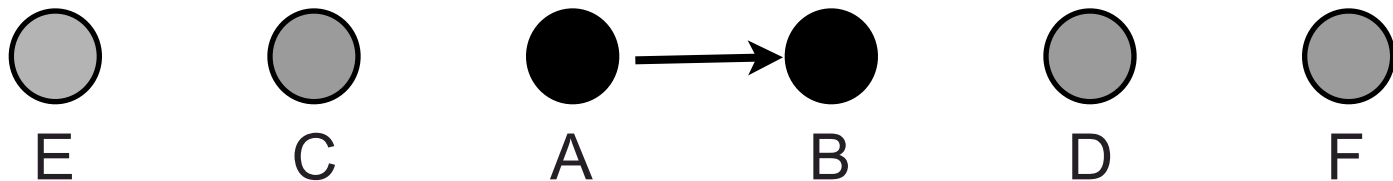
- Bluetooth low energy says that devices sleep 99% of the time
- At least the receiver must not sleep
- Parallel communication
 - Who could sleep?



A is transmitting to B

ENERGY SAVING: SLEEP!

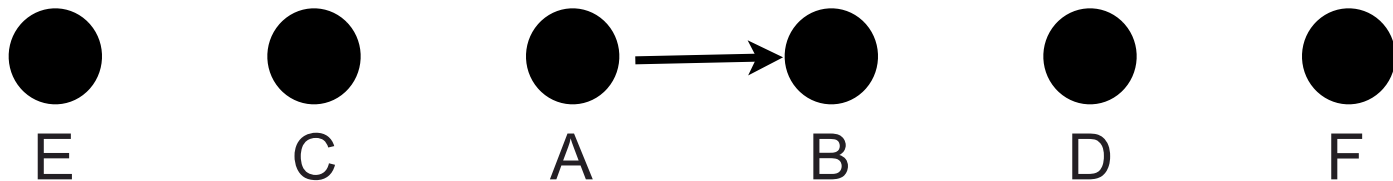
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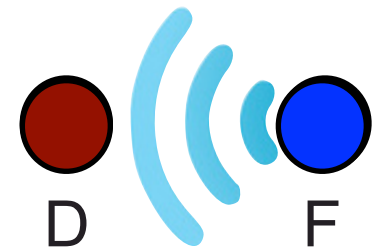
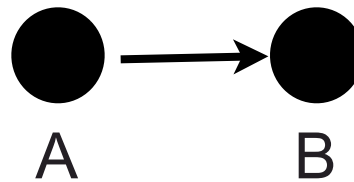
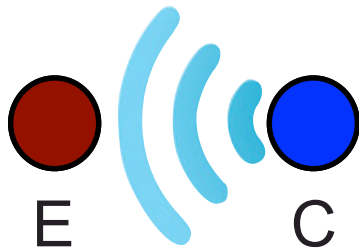
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- At least the receiver must not sleep
- Parallel communication
 - Who could sleep?
 - Who could send?



A is transmitting to B

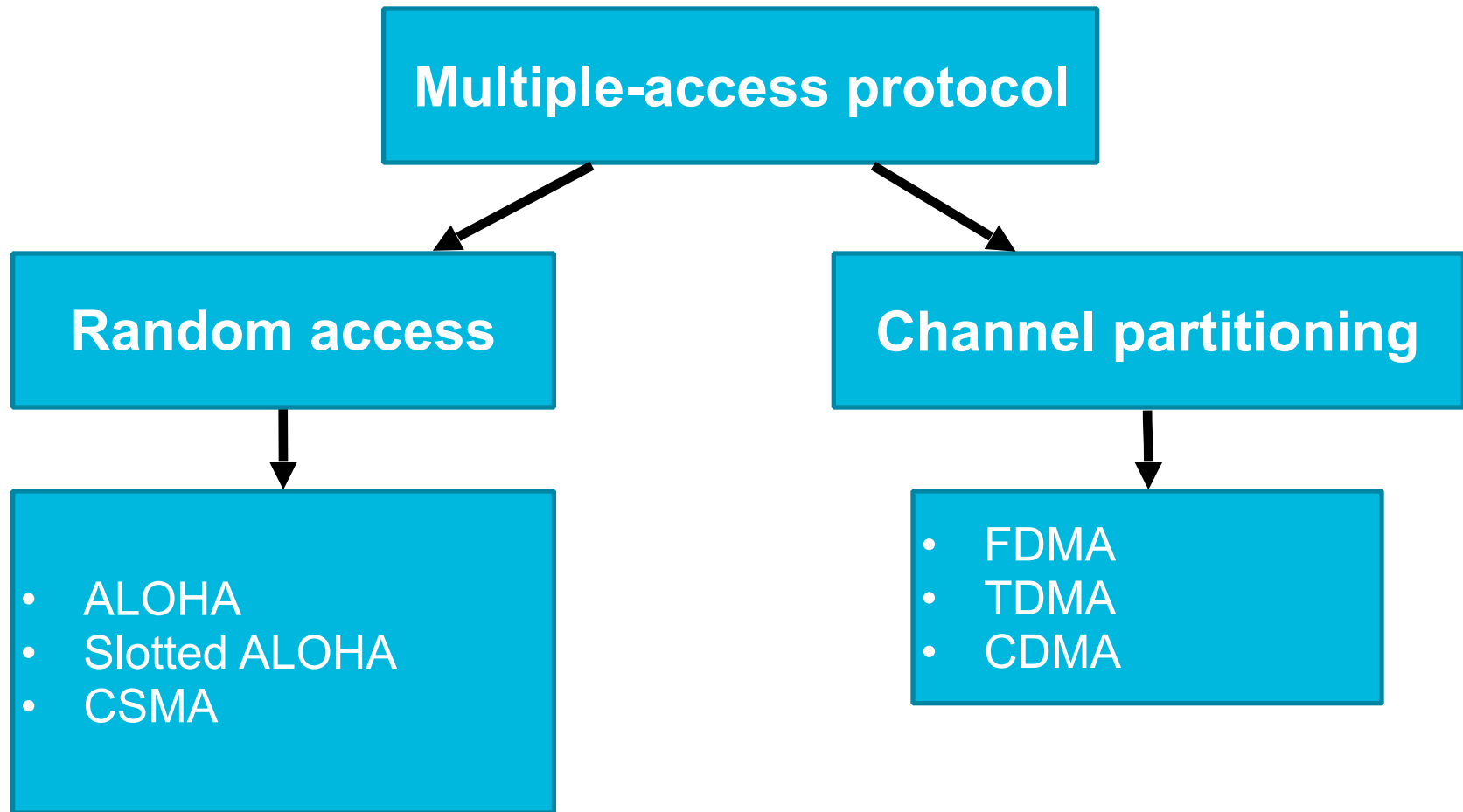
ENERGY SAVING: SLEEP!

- Bluetooth low energy says that devices sleep 99% of the time
- At least the receiver must not sleep
- Parallel communication
 - Who could sleep?
 - We can send?



A is transmitting to B

CHANNEL (MEDIUM) ACCESS METHODS



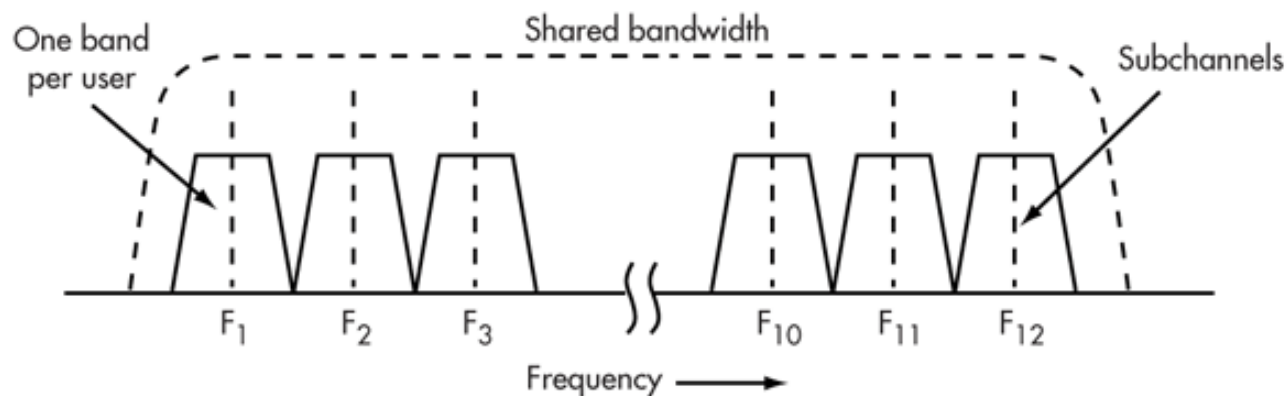
CHANNEL PARTITIONING



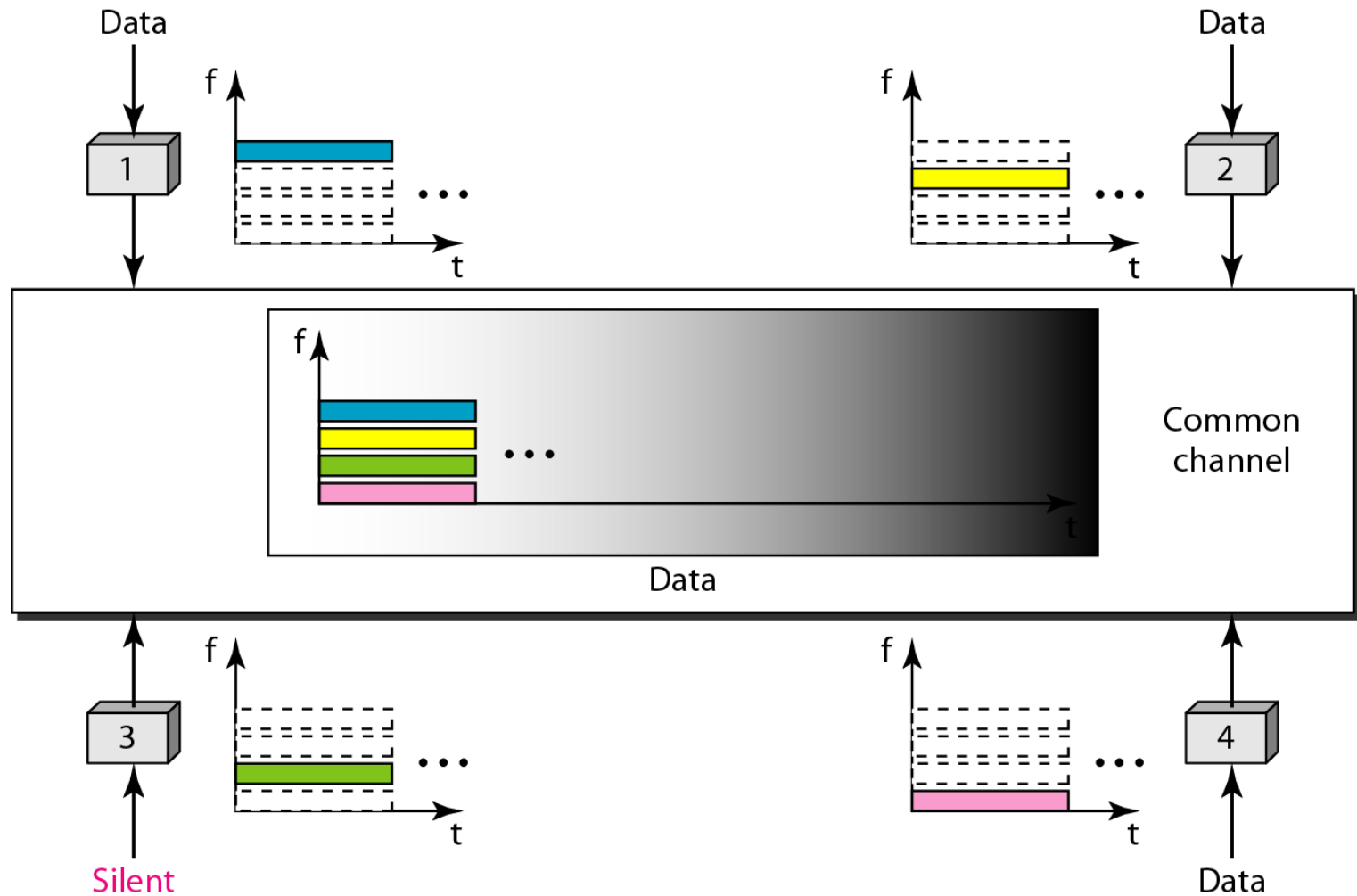
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FREQUENCY-DIVISION MULTIPLE ACCESS (FDMA) ³⁶

Under FDMA method, the available channel or bandwidth is divided into multiple individual bands that are separated by guard bands, each for use by a single user.

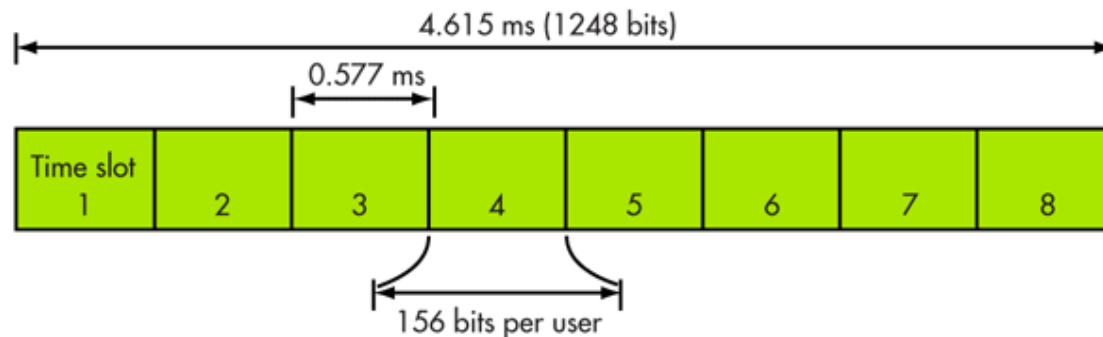


FDMA: EXAMPLE

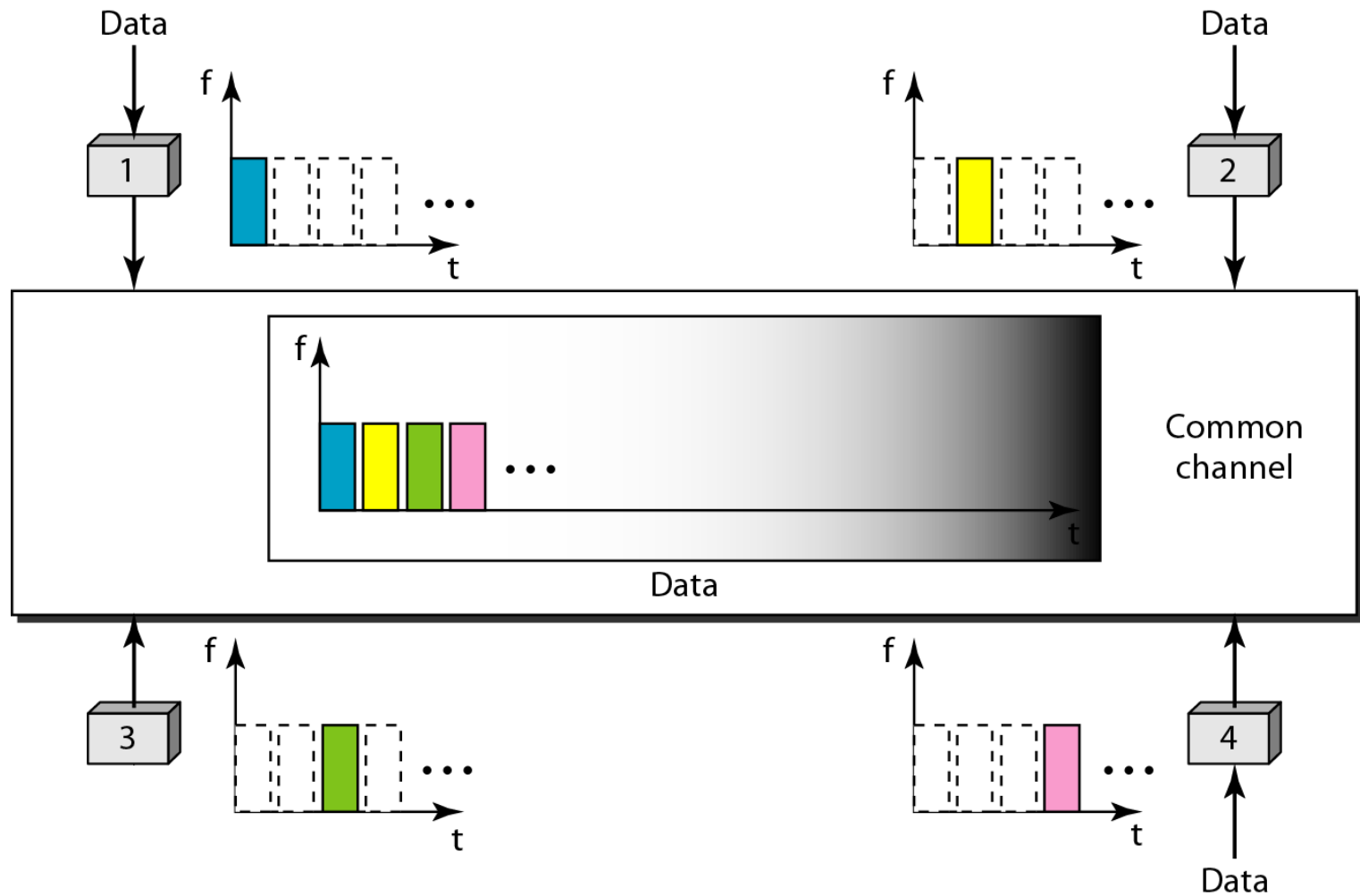


TIME-DIVISION MULTIPLE ACCESS (TDMA)

Under TDMA method, a channel is divided into time slots. Each slot is dedicated to different stations/users to transmit one byte or another digital segment.

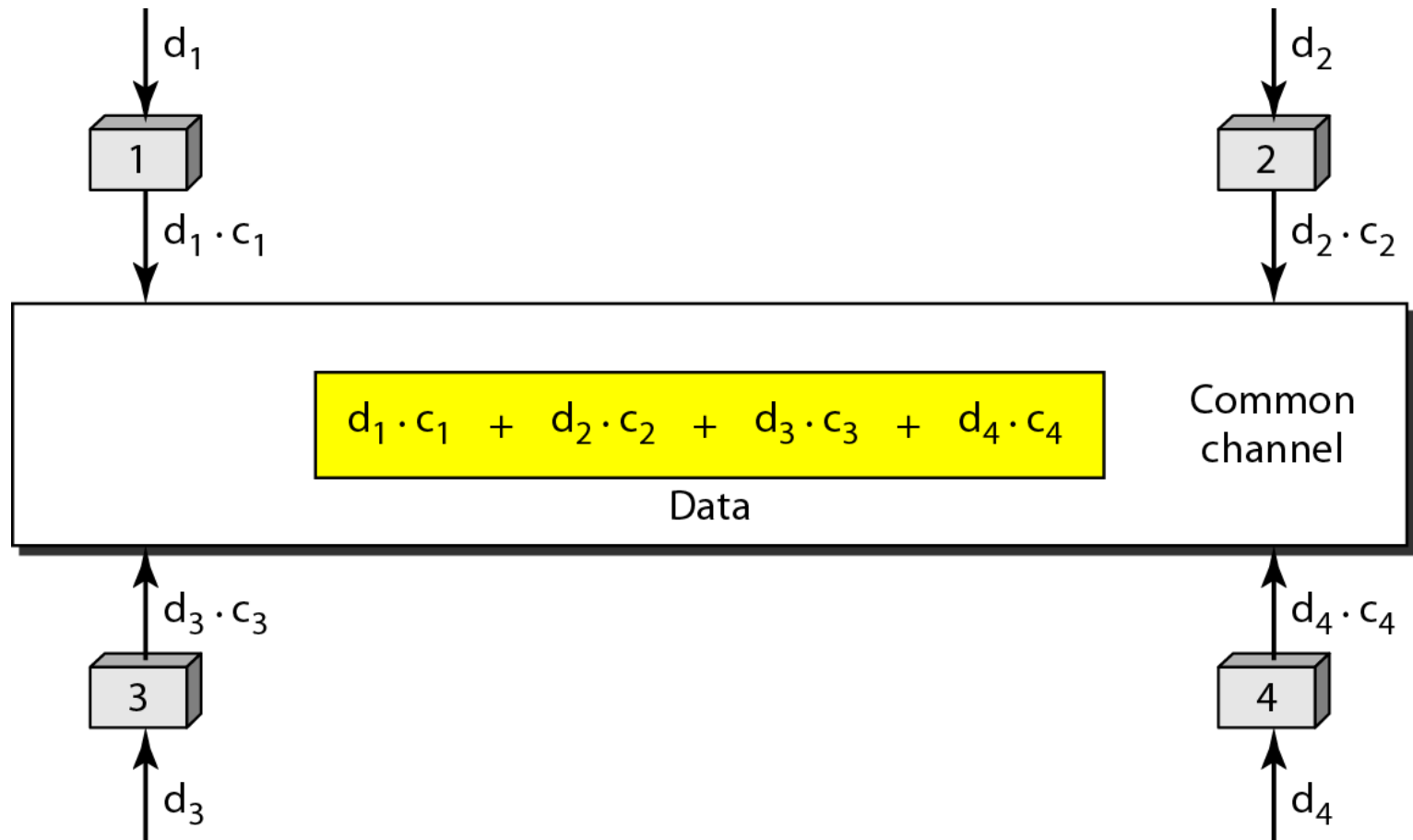


TDMA: EXAMPLE



CODE DIVISION MULTIPLE ACCESS

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RANDOM ACCESS PROTOCOLS

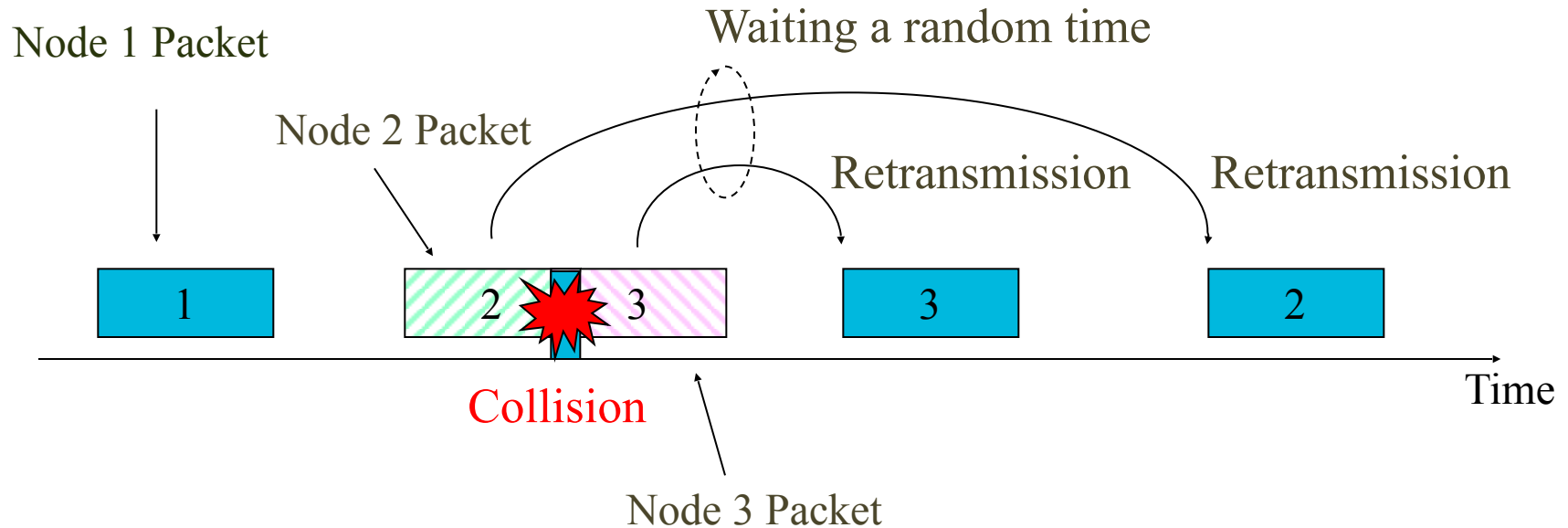


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- ▶ ALOHA: Developed in the 1970s by Hawaii University
- ▶ Whenever a node has a packet, it transmits
- ▶ Transmitter listens the **broadcast from the destination node**, to check if the transmission was successful (or not)
- ▶ Sender retransmits after some random time if there is a collision
- ▶ Works well when traffic is very limited
- ▶ Unstable when the network is loaded



ALOHA: COLLISION EXAMPLE



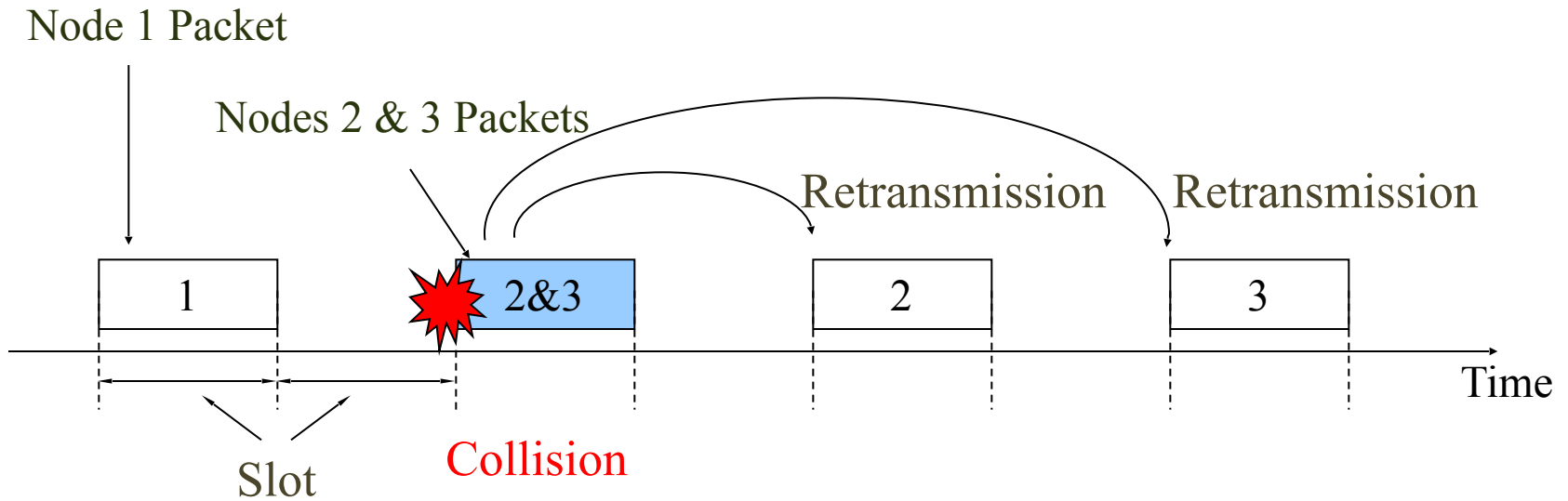
Collision mechanism in ALOHA

SLOTTED ALOHA

- ▶ Slotted ALOHA: *Time is slotted*
 - Transmission can take place only *at the beginning of slot*
 - Collision duration is reduced



SLOTTED ALOHA: COLLISION EXAMPLE

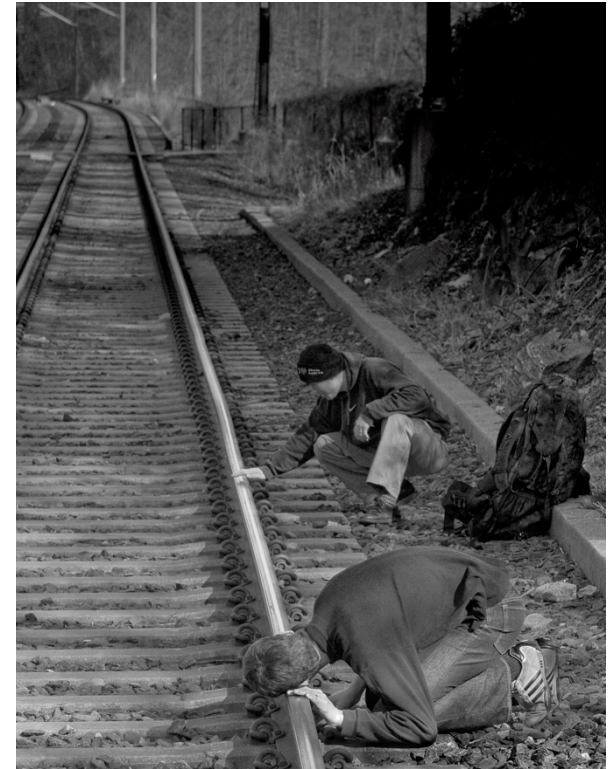


Collision mechanism in slotted ALOHA

- ▶ A node transmits **only if** no other transmission is ongoing
 - Listens to the channel before transmitting a packet

LISTENING (SAMPLING) THE MEDIUM

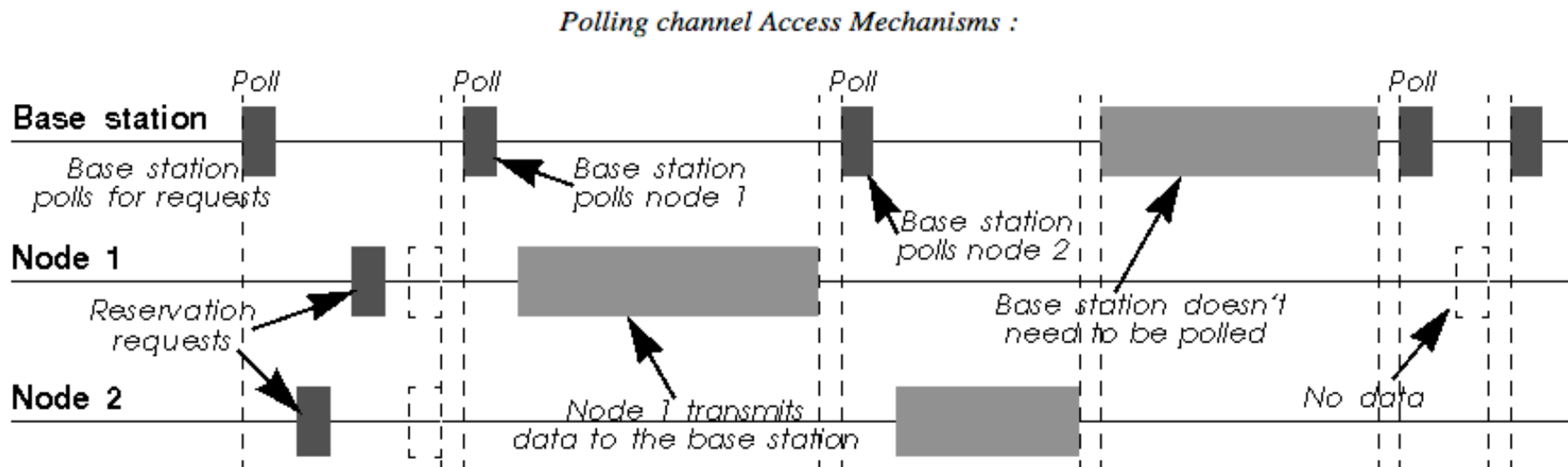
- ▶ Activity detection on the medium
- ▶ Clear Channel Assessment (CCA)
 - Detected energy $>$ threshold
 - The detected signal corresponds to the PHY layer of the node with the same modulation and propagation characteristics



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POLL

- ▶ Between TDMA and CSMA
- ▶ Base station control mechanism, with variable packets size, and a poll mechanism



NOW HOW CAN YOU SLEEP?

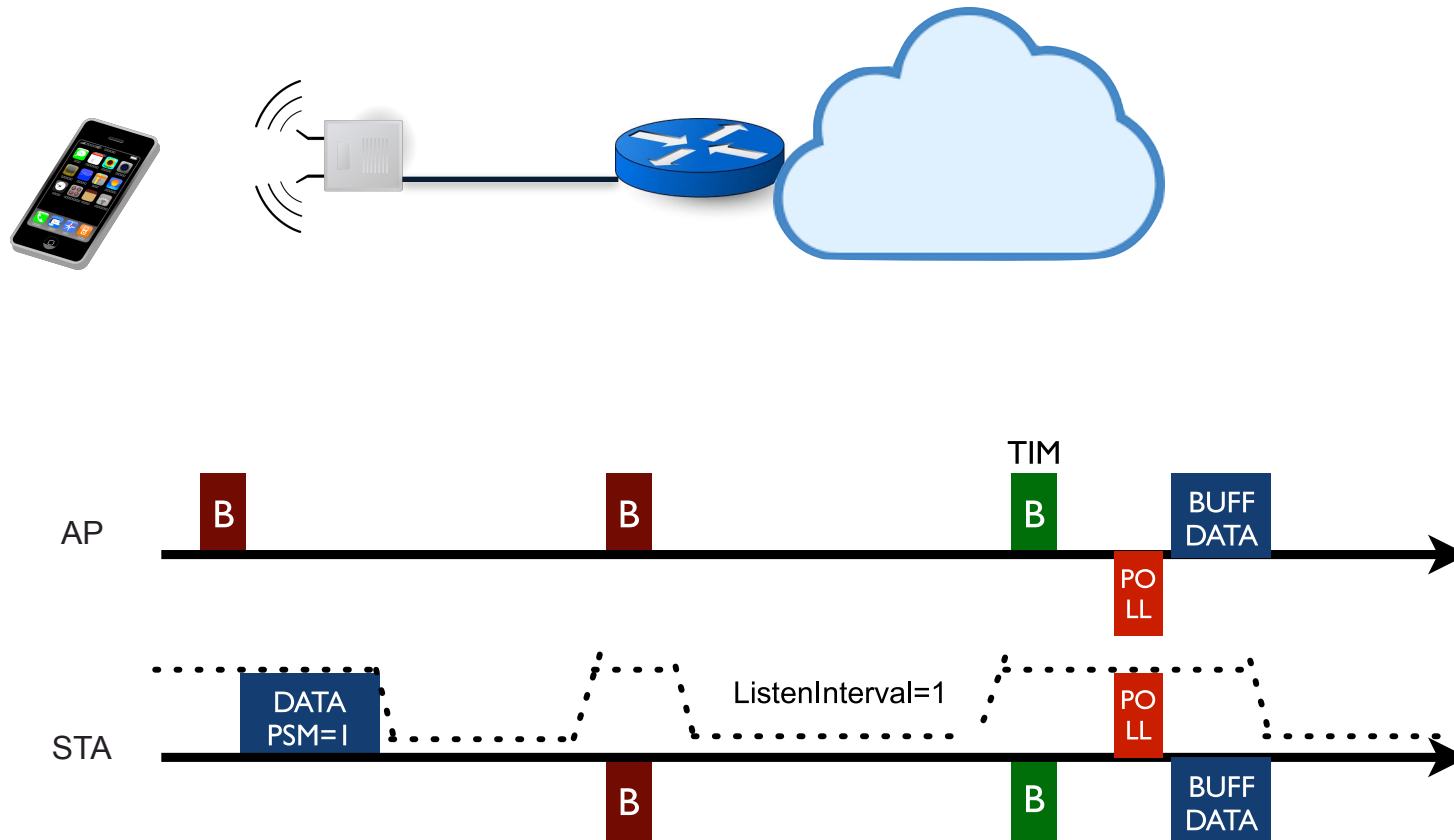


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POWER SAVING MODE (INFRASTRUCTURE)

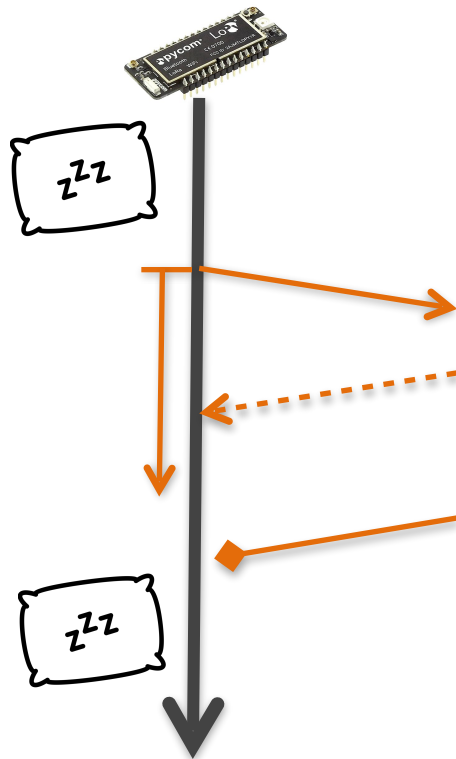


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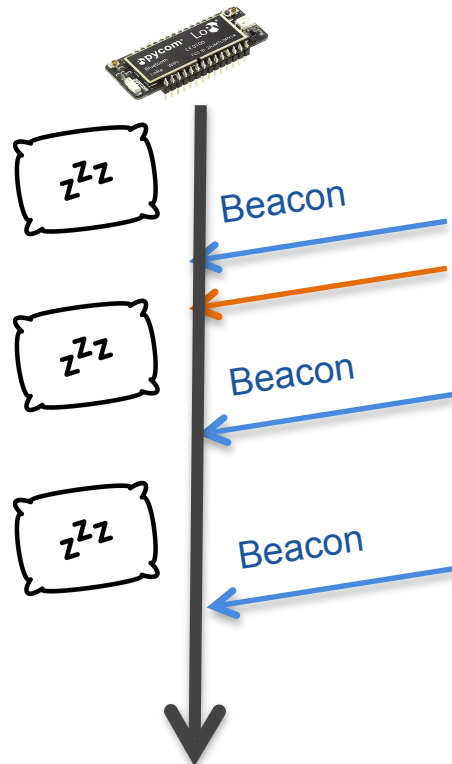


- MT activates the **PSM** bit in any frame (MAC header)
- It listens to some beacons and looks at TIM field
- It asks for buffered data (PS_POLL) and comes back to sleep

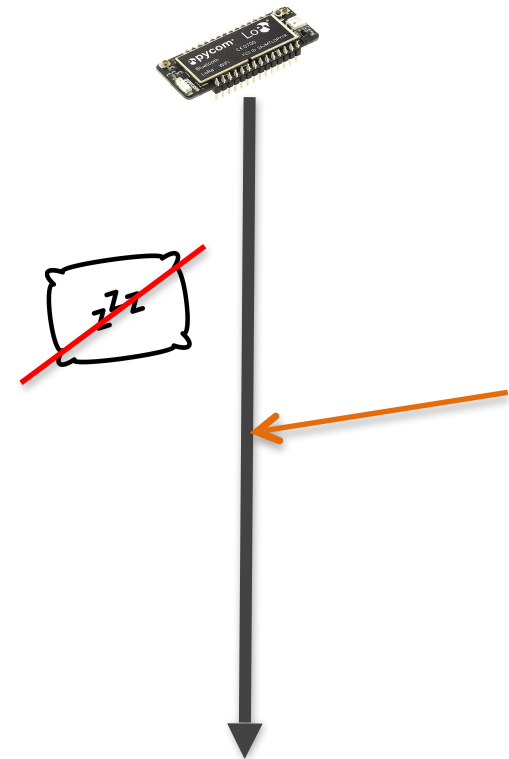
Class A

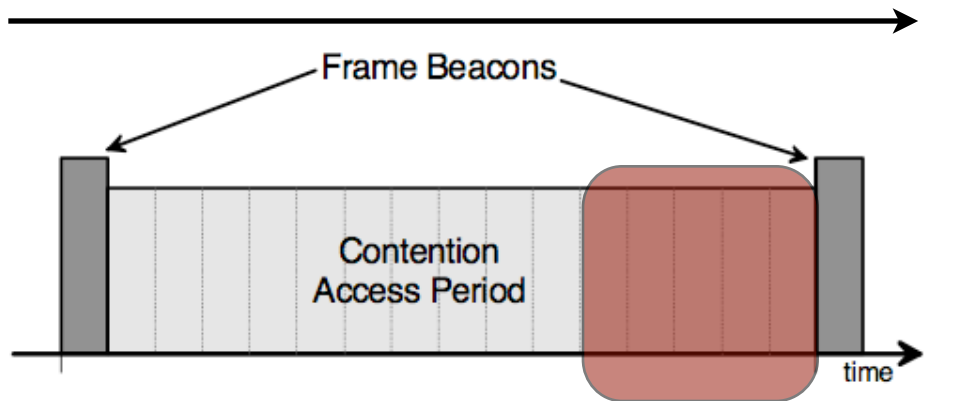


Class B



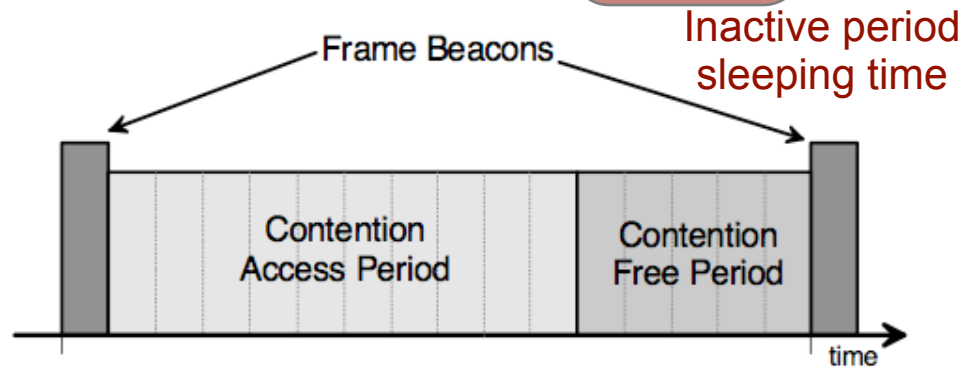
Class C



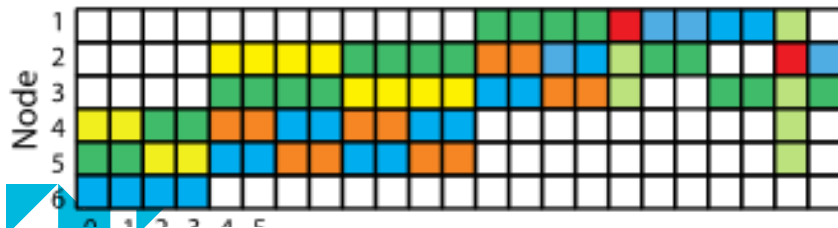


Beacon less - no synchronization

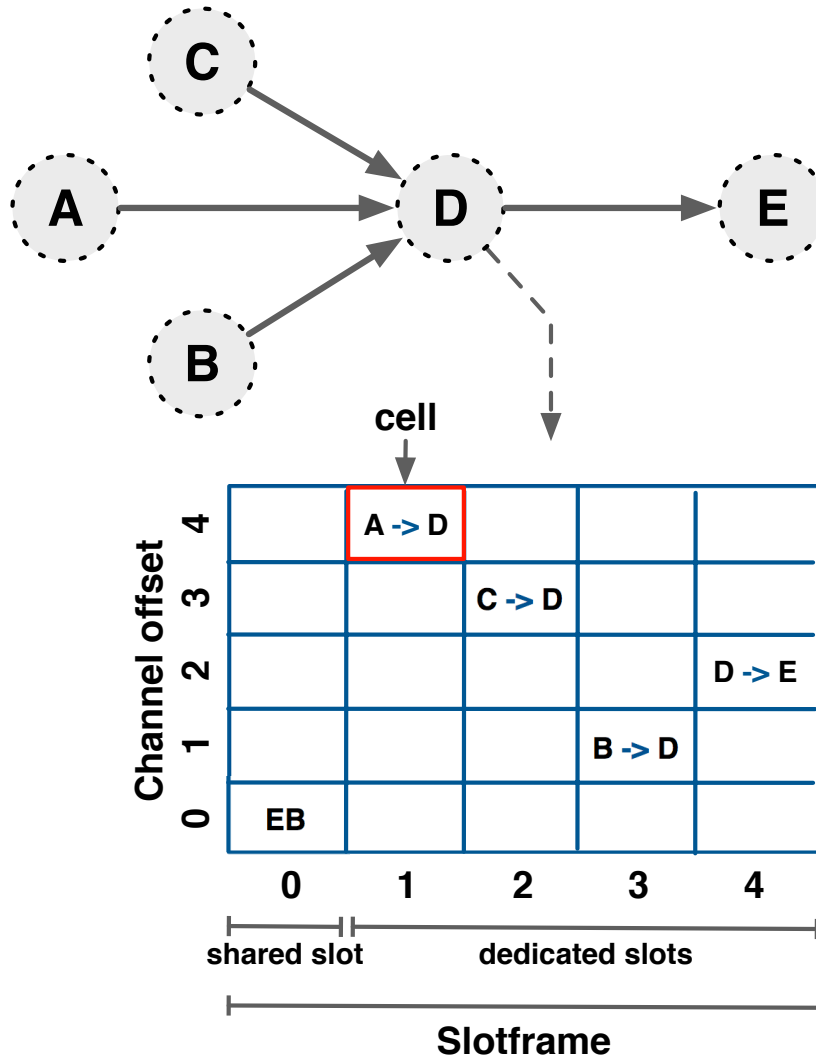
Beacon - Compete for the access (CSMA/CA)



Beacon - Competition & guaranteed service (CSMA/CA and Poll)

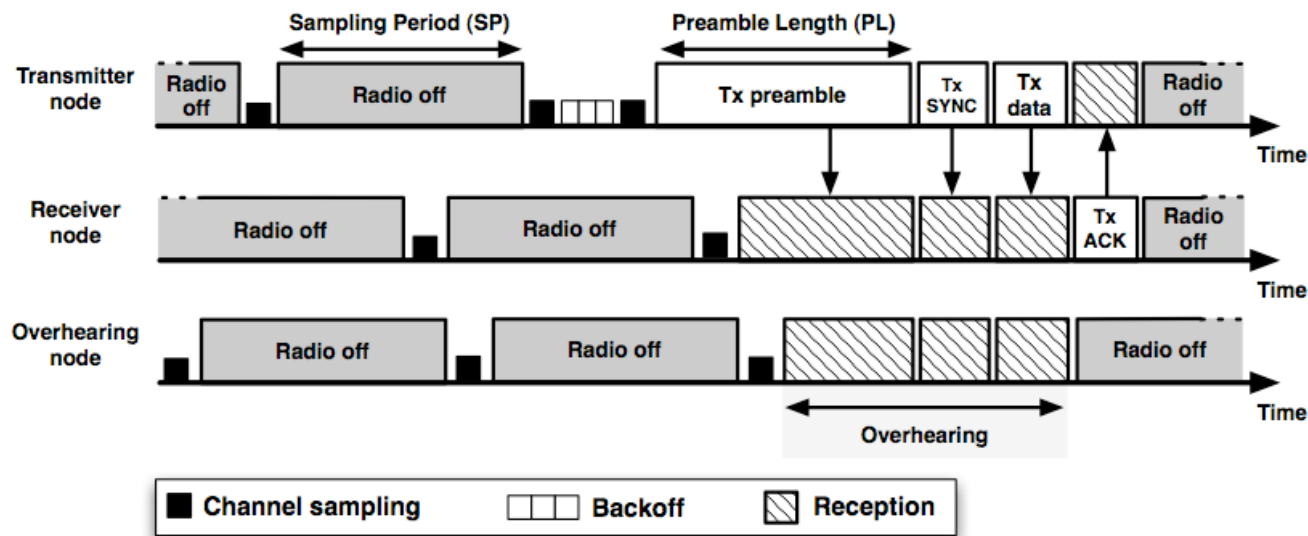


Time Slotted / Channel Hopping



ADDITIONAL MECHANISM (NON STANDARDS)

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Preamble MAC

ROUTING



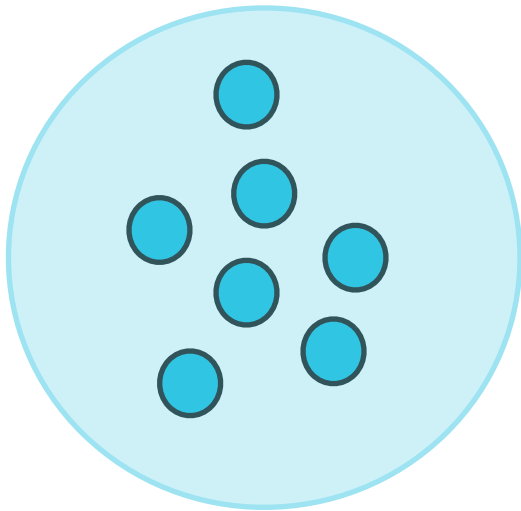
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How to go to Trinitaine?

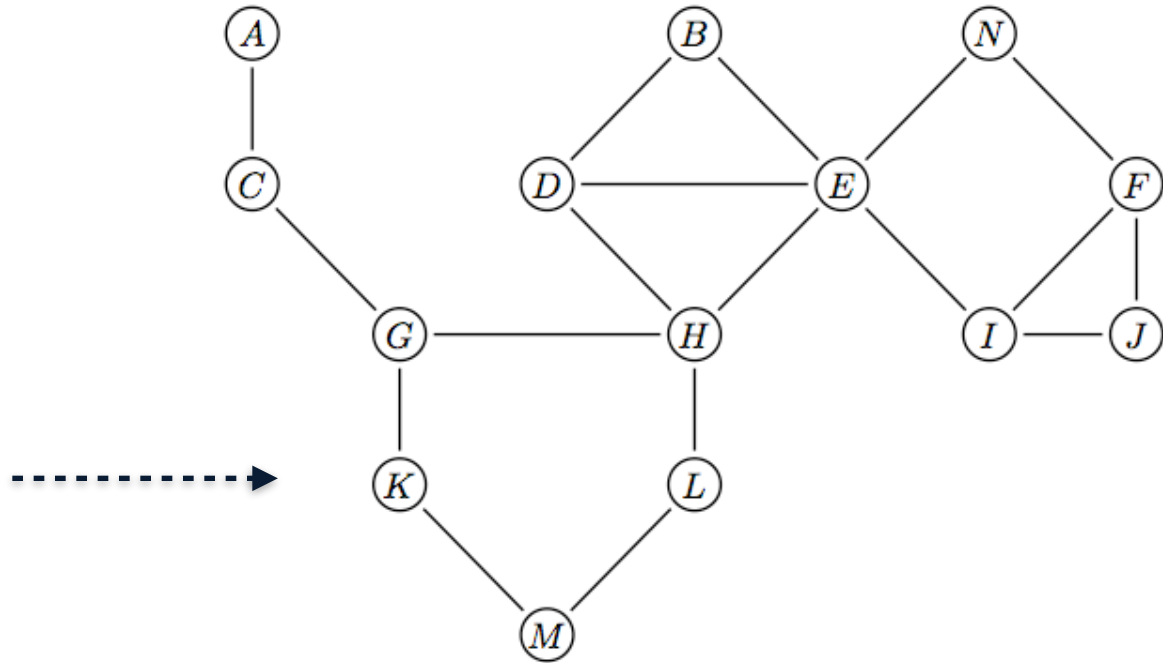


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All nodes hear each others



Extended network area

- **Routing**
 - Create and maintain the routes
 - Build a given vision of the topology
 - Objectives
 - Minimum signalling
 - No loops
 - Quick convergence
 - Take into account the link variability
 - energy consumption
- **Forwarding process**
 - Retransmission toward the destination

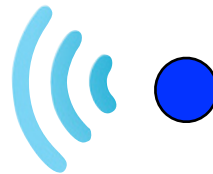
■ Neighbor discovery

- Periodically send a (one-hop) signaling message

Hello I am X

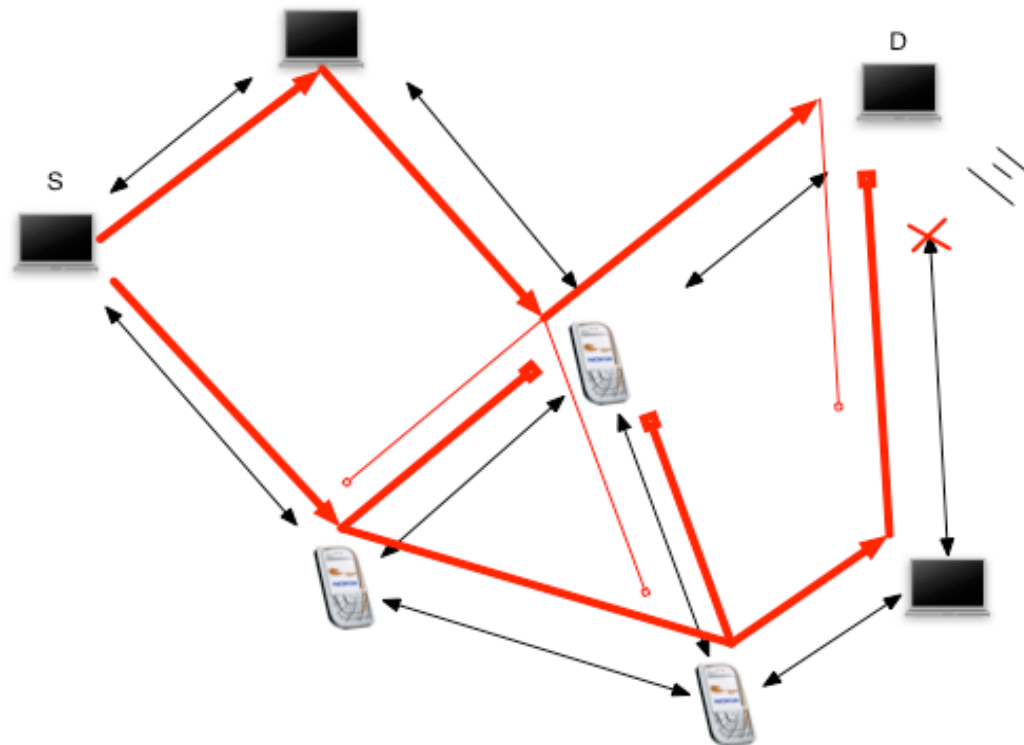


Hello I am Y



- **Flooding**

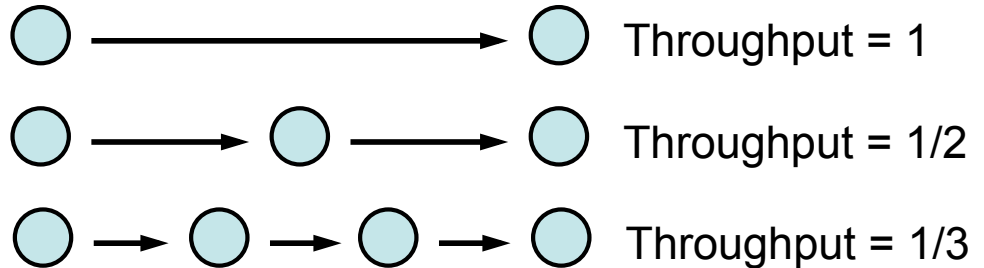
- Flood the network with a signaling information
- Learn and build the routes



■ How do you select a route / a neighbor

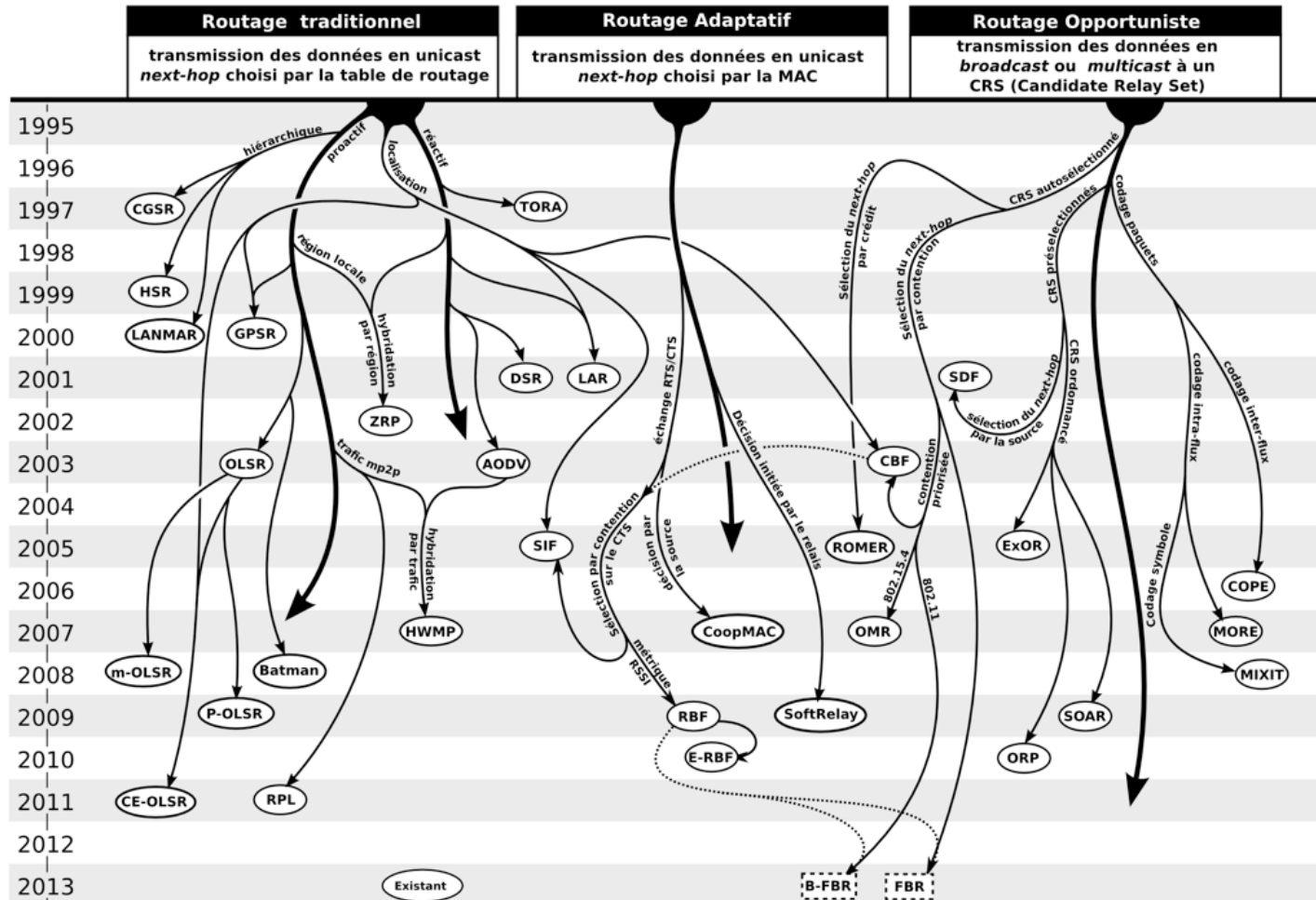
- End-to-end

- Hop count
- Packet Delivery ratio
- Delay
- throughput
- Energy consumption



- Hop-by-hop

- Expected Transmission Time (ETT)
- Expected Transmission Count (ETX)
- Energy consumption



PROACTIVE VS. REACTIVE

► Proactive

- Traditional distributed shortest-path protocols
- Based on periodic updates
- High routing overhead

► Reactive (On-demand) protocols

- Discover routes when needed
- Source-initiated route discovery

► Tradeoff

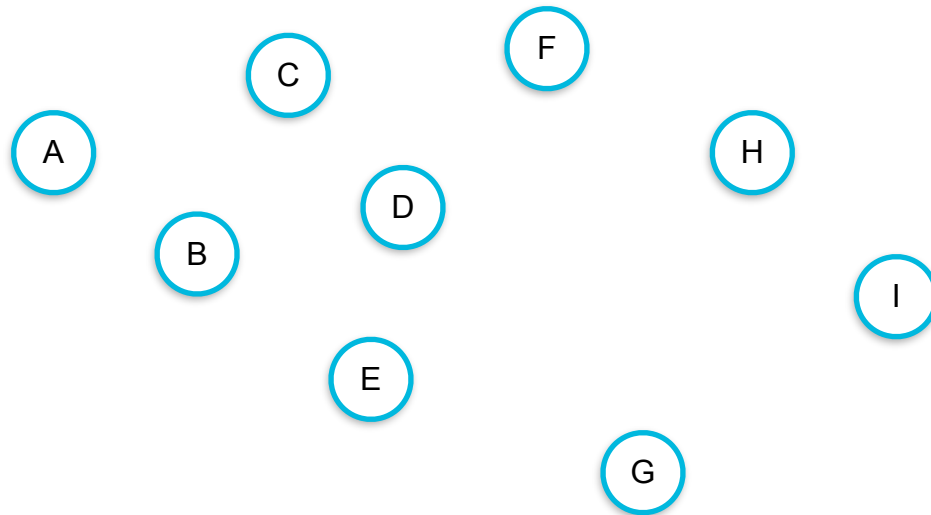
- State maintenance traffic vs. route discovery traffic
- Route via maintained route vs. delay for route discovery



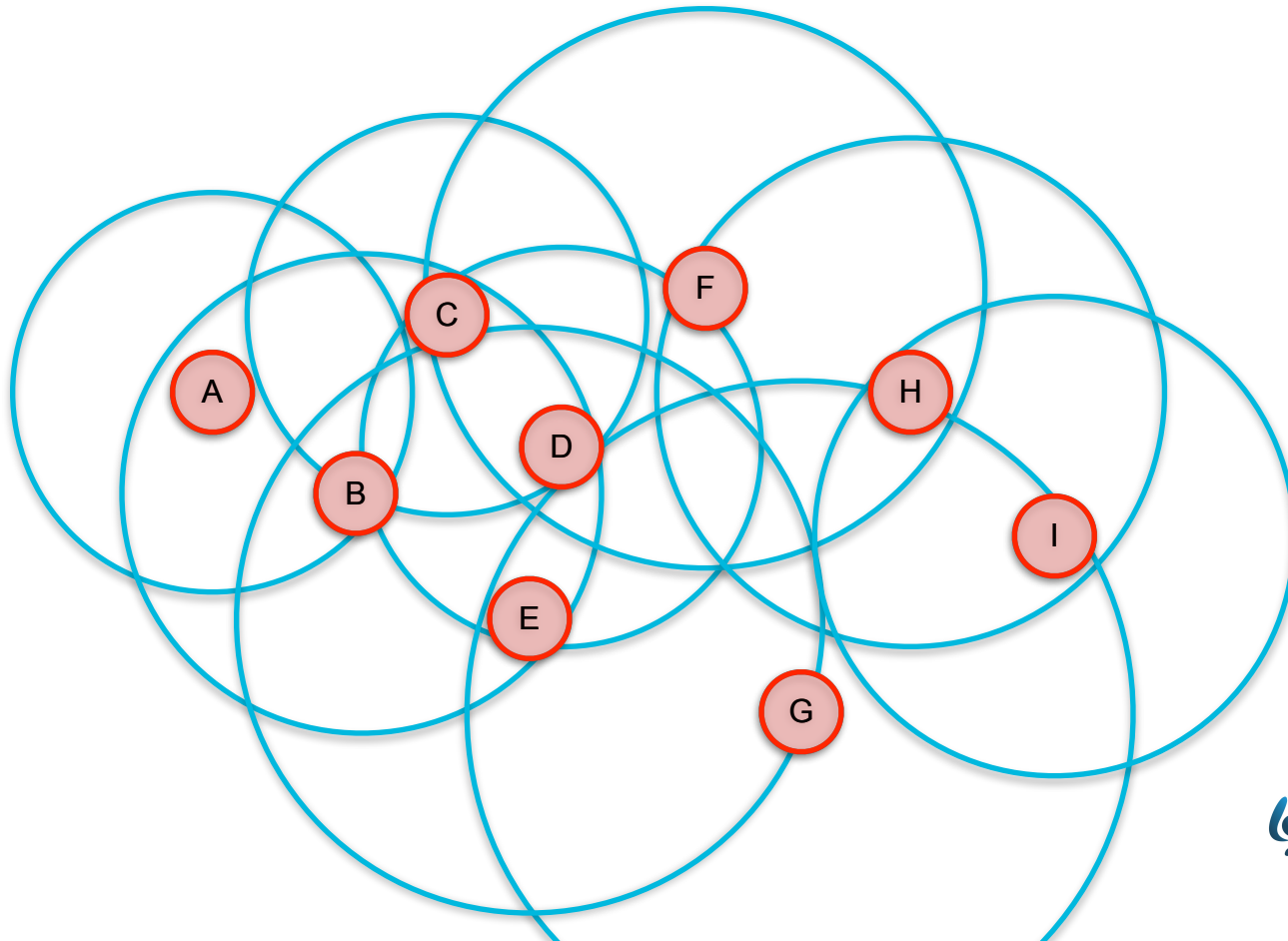
Vs.



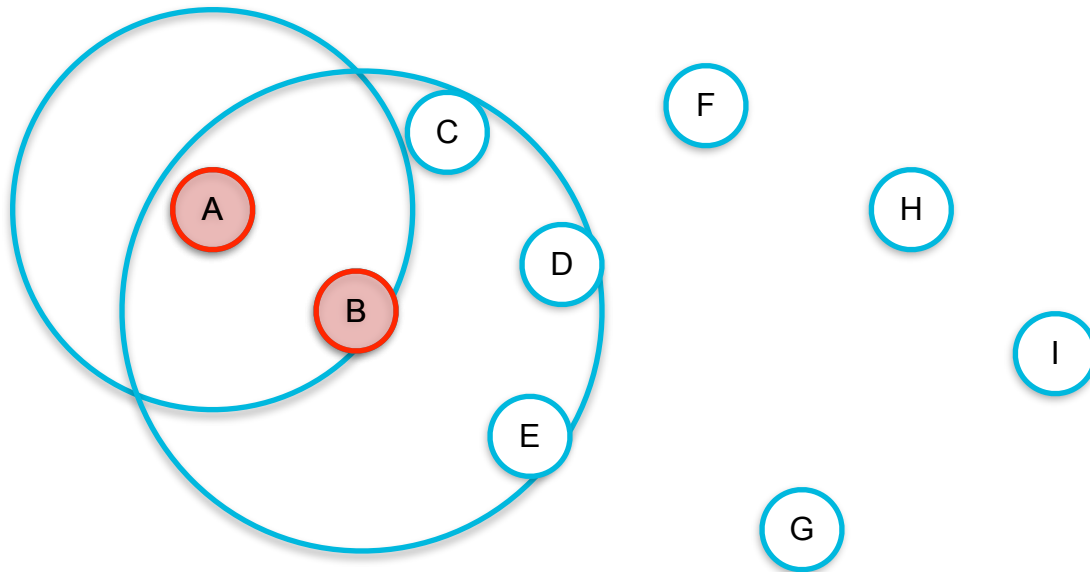
1. Flood a Route Request



1. Flood a Route Request

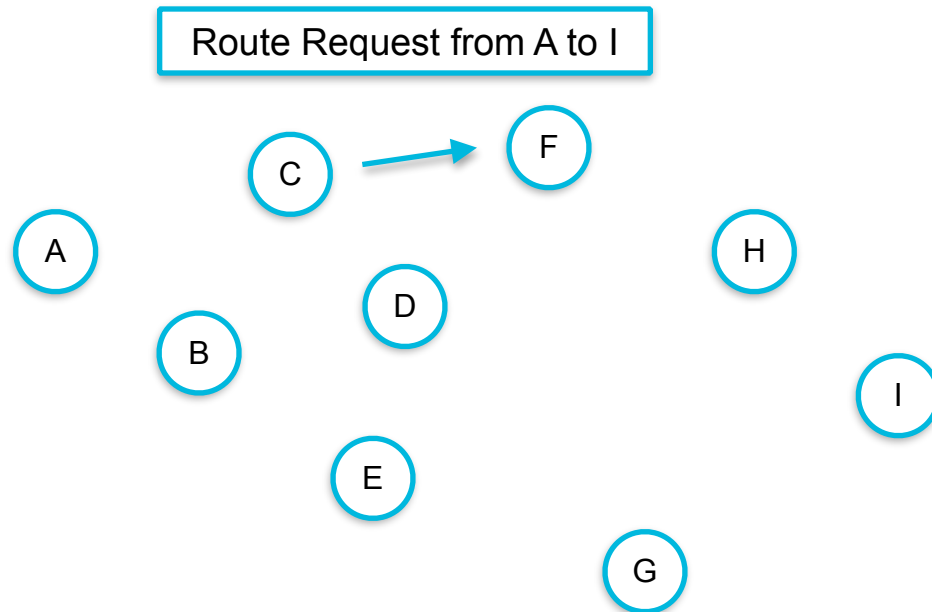


1. Flood a Route Request

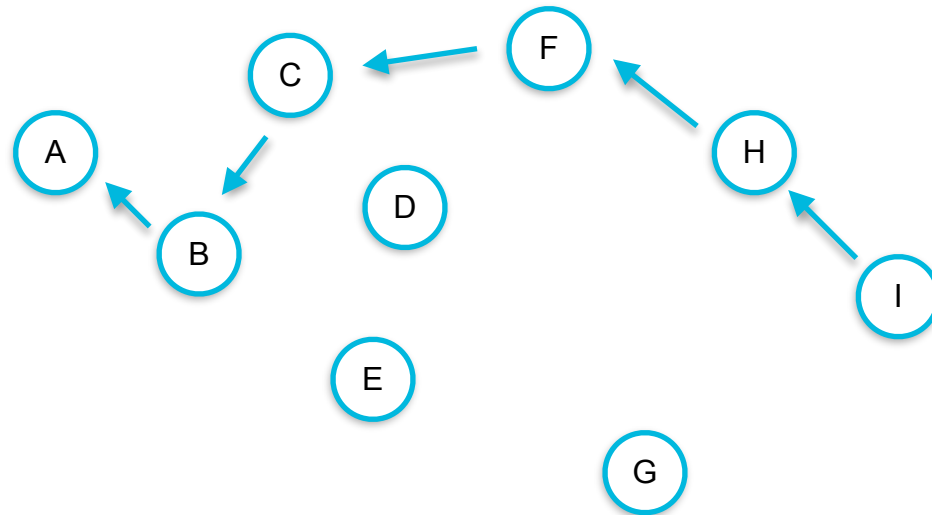


F Routing Table

Dest.	Next Hop	Dest Seq N.	Hops	Lifetime
A	C	4	3	100
C	*	10	0	100

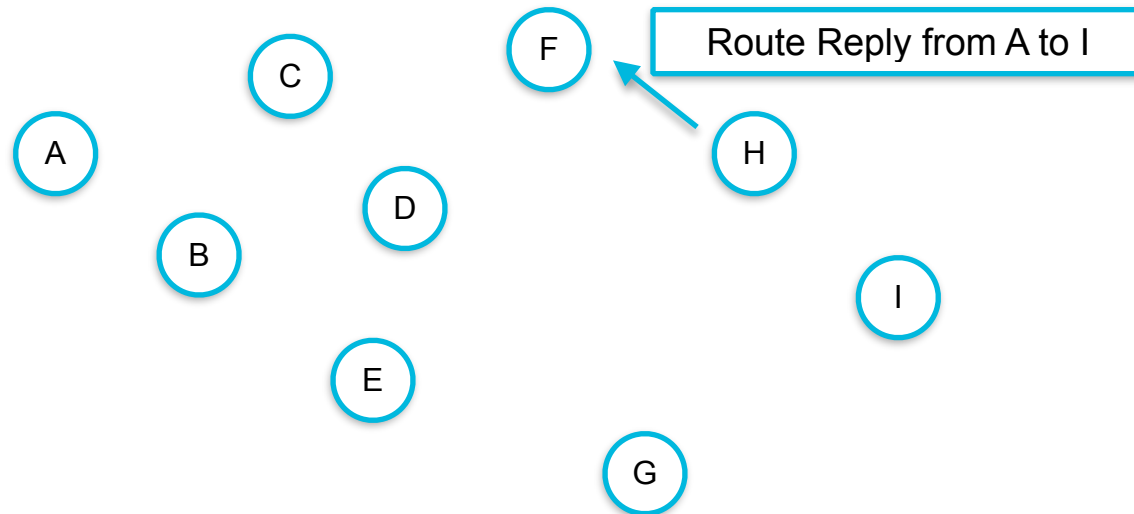


1. Flood a Route Request
2. Respond with a Route Reply

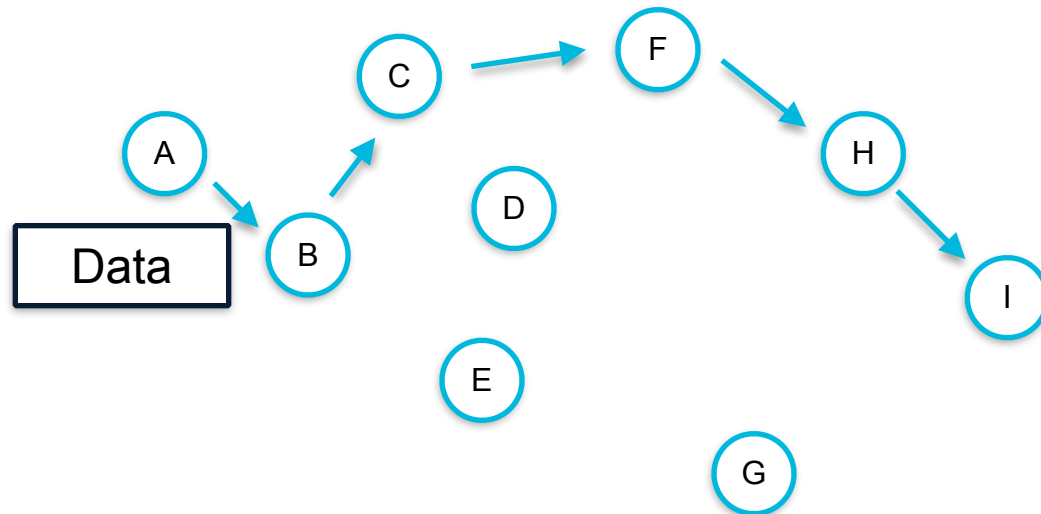



F Routing Table

Dest.	Next Hop	Dest Seq N.	Hops	Lifetime
A	C	4	3	99
C	*	10	0	99
I	H	6	2	100

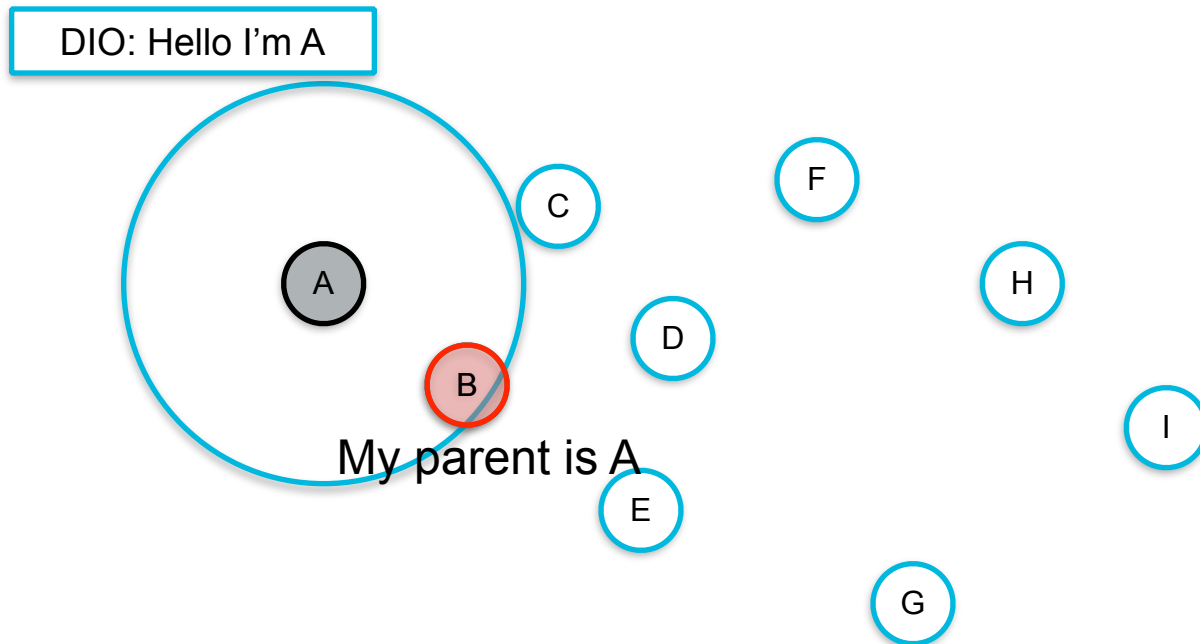


1. Flood a Route Request
2. Respond with a Route Reply
3. Send your data



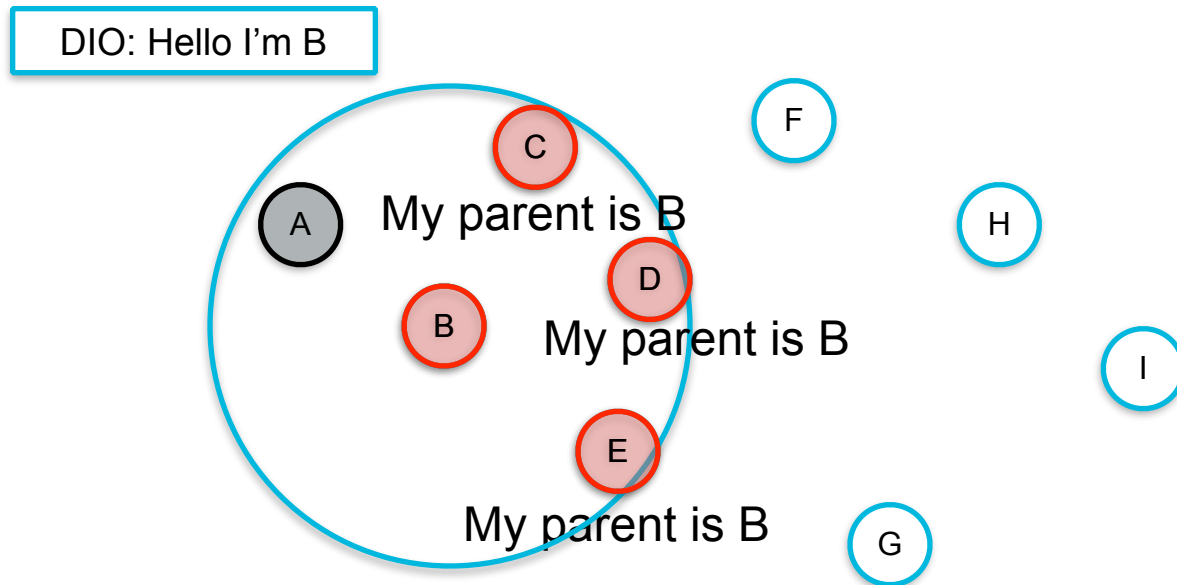
 is the root node.

1. Send DIO



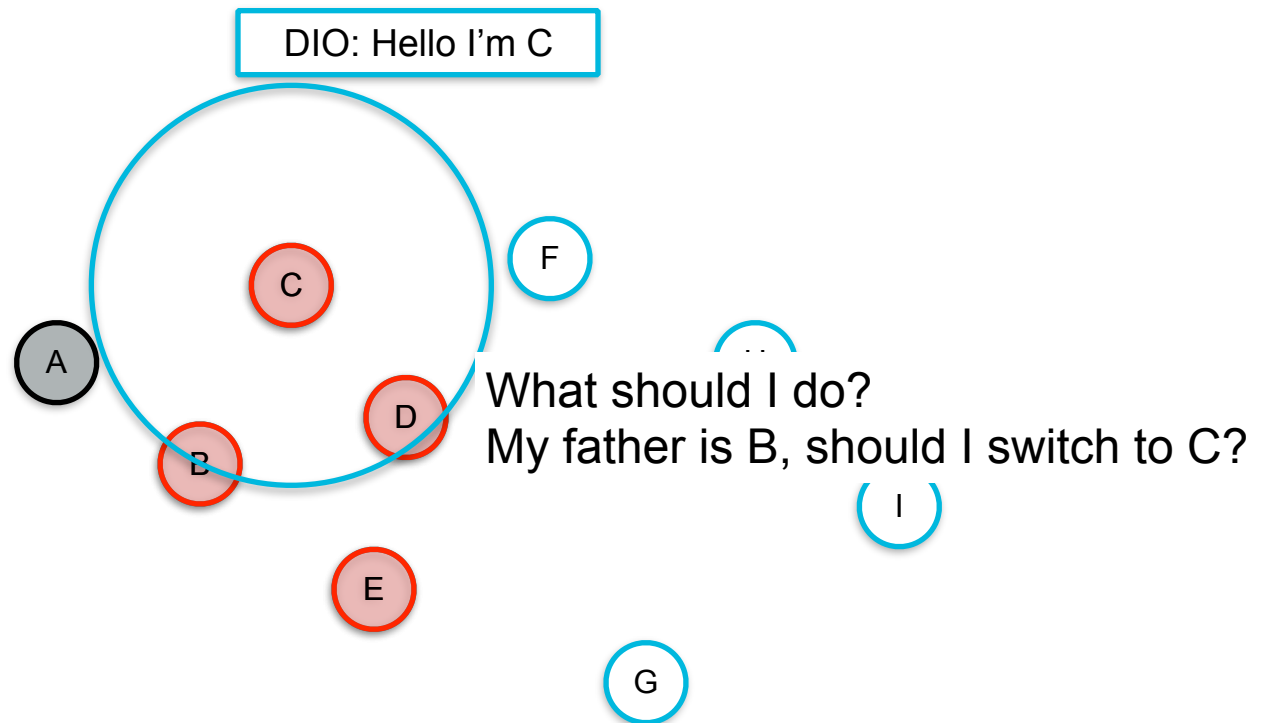
Ⓐ is the root node.

1. Send DIO



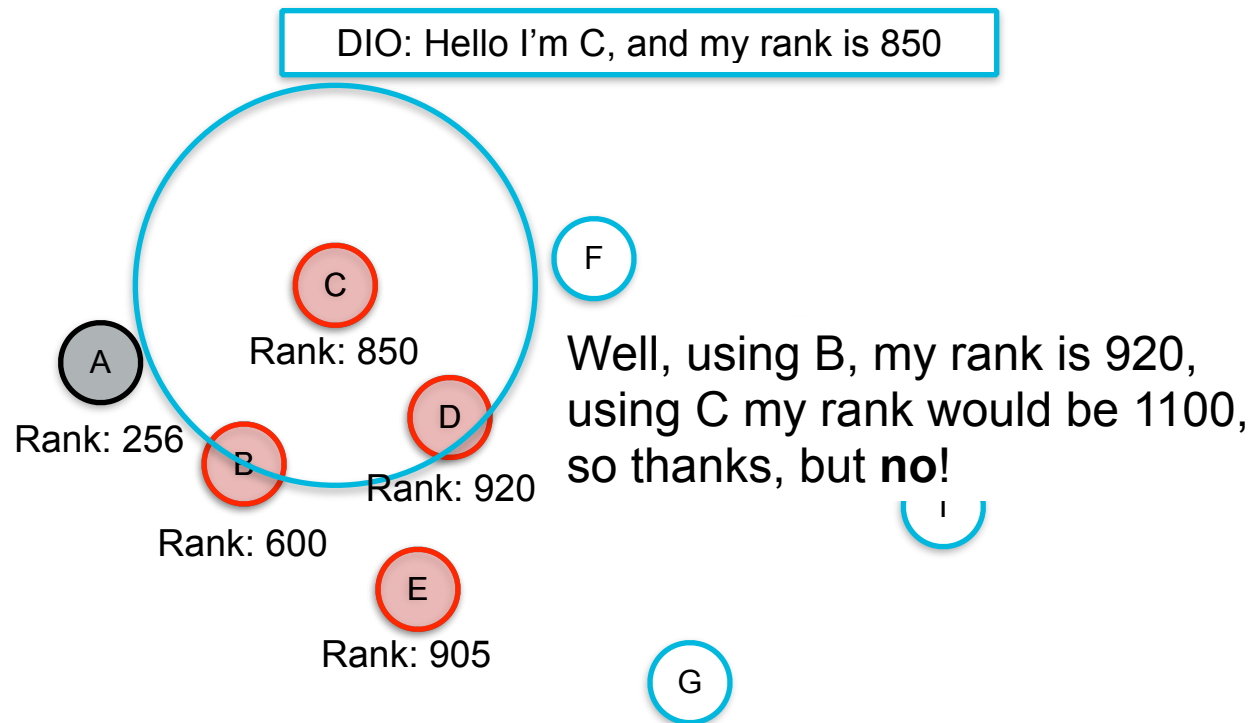
Ⓐ is the root node.

1. Send DIO



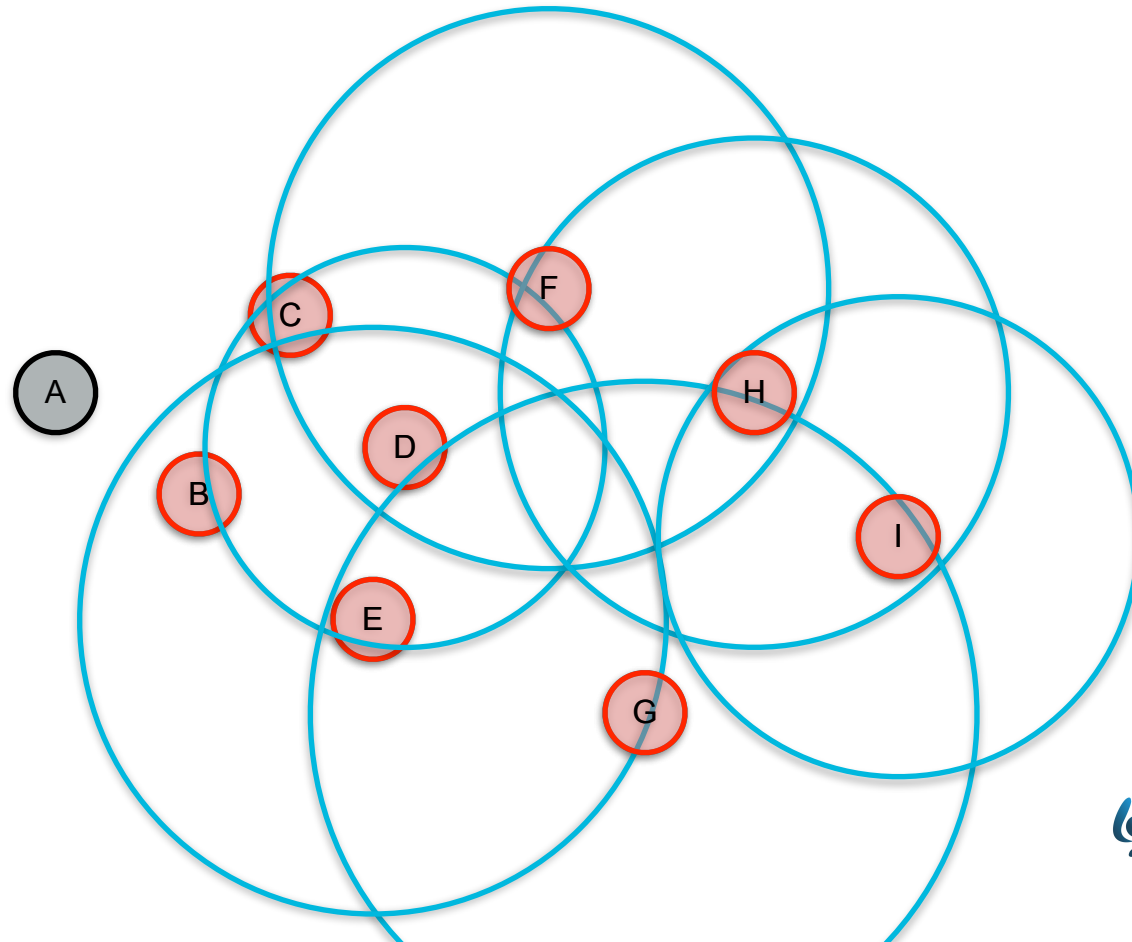
Ⓐ is the root node.

1. Send DIO



A is the root node.

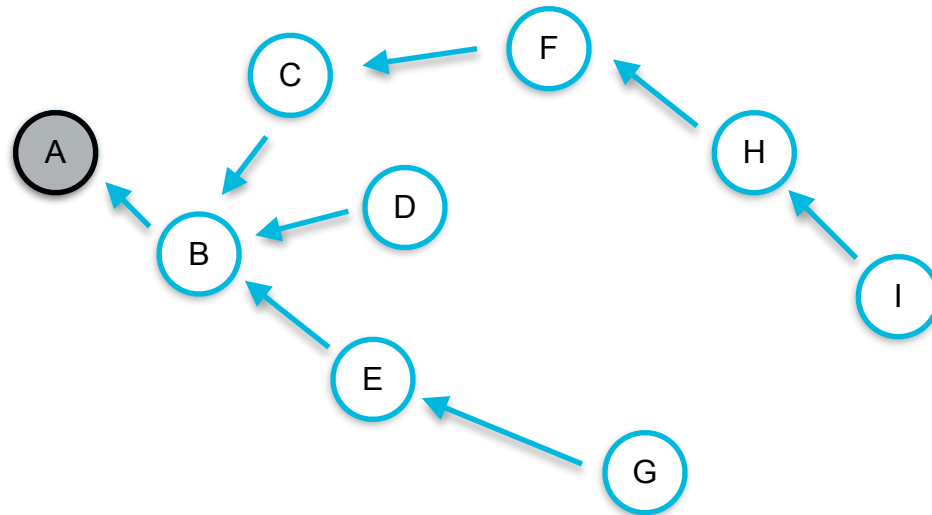
1. Send DIO



A is the root node.

1. Send DIO

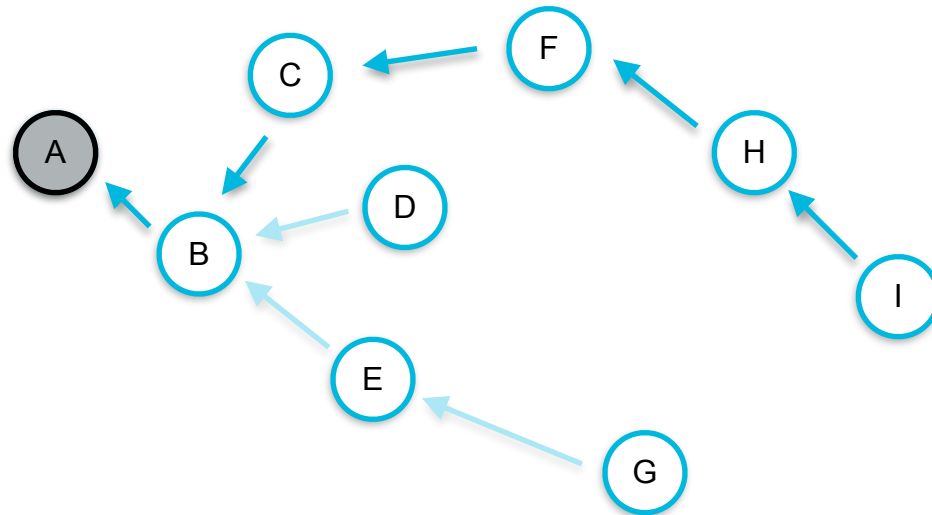
2. Build a **Destination Oriented Directed Acyclic Graph**



A is the root node.

1. Send DIO
2. Build a **Destination Oriented Directed Acyclic Graph**
3. Send data

Just send to your parent!



So which one do you prefer?

78

LOADng

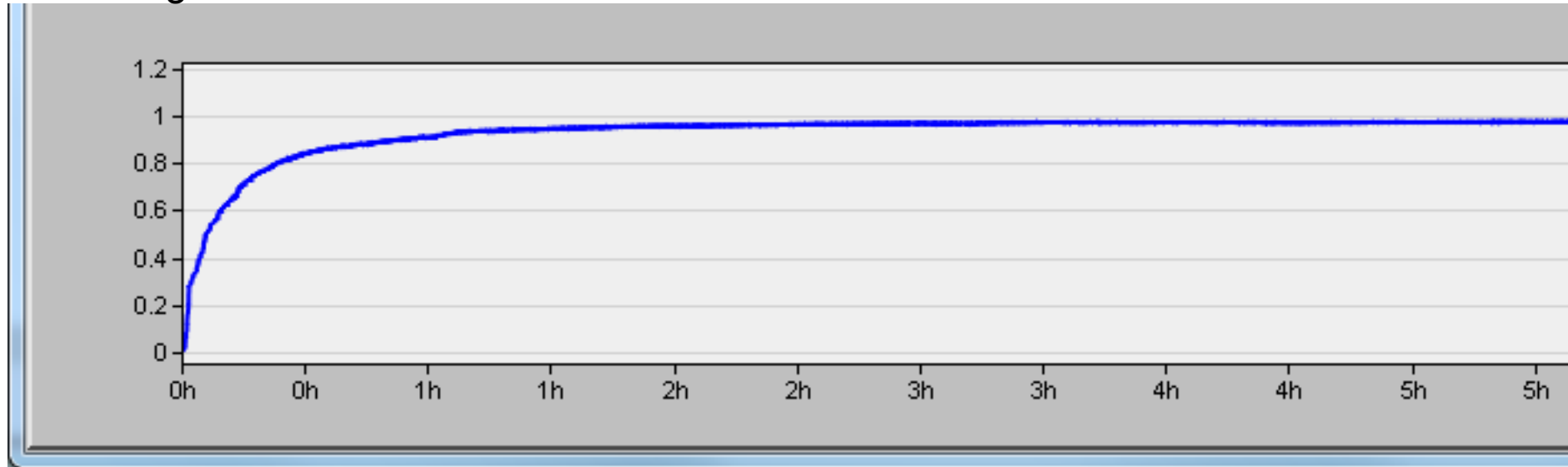
Vs

RPL

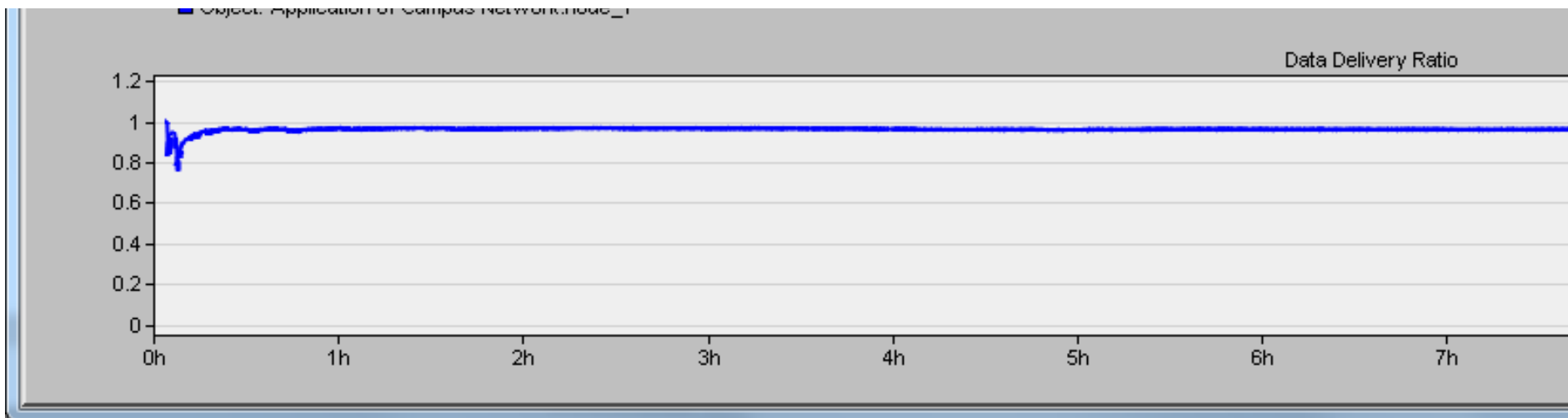
LOADng vs RPL: Packet Delivery Ratio

79

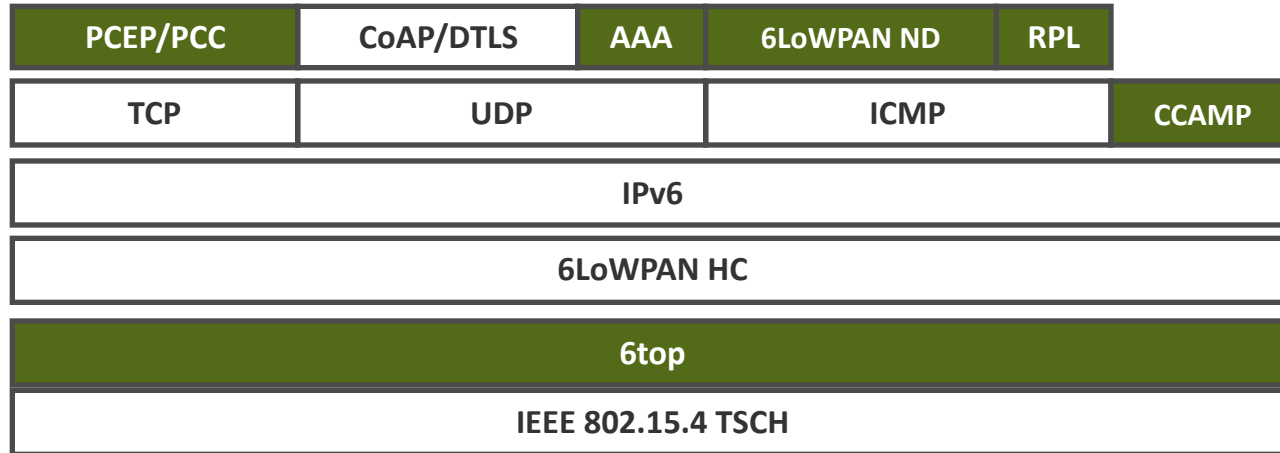
LOADng




RPL



IoT stack



Energy saving => 

Avoid verticals, but still many different applications

